arXiv:1510.00892v1 [math.LO] 4 Oct 2015

Order in Implication Zroupoids

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Abstract

The variety **I** of implication zroupoids (using a binary operation \rightarrow and a constant 0) was defined and investigated by Sankappanavar in [7], as a generalization of De Morgan algebras. Also, in [7], several new subvarieties of **I** were introduced, including the subvariety $\mathbf{I}_{2,0}$, defined by the identity: $x'' \approx x$, which plays a crucial role in this paper. Some more new subvarieties of **I** are studied in [3] that includes the subvariety **SL** of semilattices with a least element 0; and an explicit description of semisimple subvarieties of **I** is given in [5].

It is a well known fact that there is a partial order (denote it by \sqsubseteq) induced by the operation \land , both in the variety **SL** of semilattices with a least element and in the variety **DM** of De Morgan algebras. As both **SL** and **DM** are subvarieties of **I** and the definition of partial order can be expressed in terms of the implication and the constant, it is but natural to ask whether the relation \sqsubseteq on **I** is actually a partial order in some (larger) subvariety of **I** that includes both **SL** and **DM**.

The purpose of the present paper is two-fold: Firstly, a complete answer is given to the above mentioned problem. Indeed, our first main theorem shows that the variety $\mathbf{I}_{2,0}$ is a maximal subvariety of \mathbf{I} with respect to the property that the relation \sqsubseteq is a partial order on its members. In view of this result, one is then naturally led to consider the problem of determining the number of non-isomorphic algebras in $\mathbf{I}_{2,0}$ that can be defined on an *n*-element chain (herein called $\mathbf{I}_{2,0}$ -chains), *n* being a natural number. Secondly, we answer this problem in our second main theorem which says that, for each $n \in \mathbb{N}$, there are exactly *n* nonisomorphic $\mathbf{I}_{2,0}$ -chains of size *n*.

l Introduction

The widely known fact that Boolean algebras can be defined using only implication and a constant was extended to De Morgan algebras in [7]. The crucial role played by a certain identity, called (I), led Sankappanavar, in 2012, to define and investigate, the variety I of implication zroupoids (I-zroupoids) generalizing De Morgan algebras. Also, in [7], he introduced several new subvarieties of I and found some relationships among those subvarieties. One of the subvarieties of I, called $I_{2,0}$, defined by the identity: $x'' \approx x$ and studied in [7], plays a crucial role in this paper. In [3], we introduce several more new subvarieties of I, including the subvariety SL which is term-equivalent to the (well known) variety of \lor -semilattices with a least element 0, and describe further relationships among the subvarieties of I. An explicit description of semisimple subvarieties of I is given in [5].

It is also a well known fact that there is a partial order induced by the operation \wedge , both in the variety **SL** of semilattices with a least element and in the variety **DM** of De Morgan algebras. As both **SL** and **DM** are subvarieties of **I** and the definition of partial order can be expressed in terms of the implication and constant, it is but natural to ask whether the relation \sqsubseteq (now defined) on **I** is actually a partial order in some (larger) subvariety of **I** that includes both **SL** and **DM**. The purpose of the present paper is two-fold: Firstly, a complete answer is given to the above mentioned problem. Indeed, our first main theorem shows that the variety $\mathbf{I}_{2,0}$ is a maximal subvariety of \mathbf{I} with respect to the property that the relation \sqsubseteq , defined by:

 $x \sqsubseteq y$ if and only if $(x \to y')' = x$, for $x, y \in \mathbf{A}$ and $\mathbf{A} \in \mathbf{I}$,

is a partial order. In view of this result, one is then naturally led to consider the problem of determining the number of non-isomorphic algebras in $\mathbf{I}_{2,0}$ ($\mathbf{I}_{2,0}$ -chains) that can be defined on an *n*-element set, *n* being a nutural number. Secondly, we answer this problem in our second main result which says that, for each $n \in \mathbb{N}$, there are exactly *n* nonisomorphic $\mathbf{I}_{2,0}$ -chains of size *n*.

2 Preliminaries

In this section we recall some definitions and results from [3], [5] and [7] that will be needed for this paper. Basic references are [1] and [2].

Definition 2.1 [7] A groupoid with zero (*zroupoid*, for short) is an algebra $\mathbf{A} = \langle A, \to, 0 \rangle$, where \to is a binary operation and 0 is a constant. A zroupoid $\mathbf{A} = \langle A, \to, 0 \rangle$ is an *implication zroupoid* (I-zroupoid, for short) if the following identities hold in \mathbf{A} , where $x' := x \to 0$:

(I)
$$(x \to y) \to z \approx [(z' \to x) \to (y \to z)']'$$

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(I_0) \quad 0'' \approx 0.
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The variety of I-zroupoids is denoted by I.

In this paper we use the characterizations of De Morgan algebras, Kleene algebras and Boolean algebras (see [7]), and semilattices with least element 0 (see [3]), as definitions.

Definition 2.2 An implication zroupoid $\mathbf{A} = \langle A, \rightarrow, 0 \rangle$ is a De Morgan algebra (**DM**-algebra, for short) if **A** satisfies the axiom:

(DM) $(x \to y) \to x \approx x$.

A **DM**-algebra $\mathbf{A} = \langle A, \rightarrow, 0 \rangle$ is a Kleene algebra (**KL**-algebra, for short) if **A** satisfies the axiom:

- (KL₁) $(x \to x) \to (y \to y)' \approx x \to x$
 - or, equivalently,
- (KL₂) $(y \to y) \to (x \to x) \approx x \to x$.

A **DM**-algebra $\mathbf{A} = \langle A, \rightarrow, 0 \rangle$ is a Boolean algebra (**BA**-algebra, for short) if **A** satisfies the axiom:

(BA) $x \to x \approx 0'$.

An implication zroupoid $\mathbf{A} = \langle A, \rightarrow, 0 \rangle$ is a semilattice with 0 (SL-algebra, for short) if \mathbf{A} satisfies the axioms:

(SM1) $x' \approx x$

(SM2) $x \to y \approx y \to x$. (Commutativity).

We denote by DM, KL, BA and SL, respectively, the variety of DM-algebras, KL-algebras, BA-algebras, and SL-algebras.

We recall from [7] the definition of another subvariety of \mathbf{I} , namely $\mathbf{I}_{2,0}$, which plays a fundamental role in this paper.

Definition 2.3 $I_{2,0}$ denotes the subvariety of I defined by the identity:

 $x'' \approx x.$

We note that **DM**, **KL**, **BA** and **SL** are all subvarieties of $I_{2,0}$ (see [7] and [3]).

Lemma 2.4 [7, Theorem 8.15] Let A be an I-zroupoid. Then the following are equivalent:

- (a) $0' \to x \approx x$
- (b) $x'' \approx x$
- (c) $(x \to x')' \approx x$
- (d) $x' \to x \approx x$.

Lemma 2.5 [7] Let $\mathbf{A} \in \mathbf{I}_{2,0}$. Then

- (a) $x' \to 0' \approx 0 \to x$
- (b) $0 \to x' \approx x \to 0'$.

Several identities true in $\mathbf{I}_{2,0}$ are given in [3], [5] and [7]. Some of those that are needed for this paper are listed in the next lemma, which also presents some new identities of $\mathbf{I}_{2,0}$ that will be useful later in this paper. The proof of the lemma is given in the Appendix.

Lemma 2.6 Let $\mathbf{A} \in \mathbf{I}_{2,0}$. Then \mathbf{A} satisfies:

(1)
$$(x \to 0') \to y \approx (x \to y') \to y$$

(2) $(0 \to x') \to (y \to x) \approx y \to x$
(3) $(y \to x)' \approx (0 \to x) \to (y \to x)'$
(4) $[x \to (y \to x)']' \approx (x \to y) \to x$
(5) $(y \to x) \to y \approx (0 \to x) \to y$
(6) $0 \to x \approx 0 \to (0 \to x)$
(7) $0 \to [(0 \to x) \to (0 \to y')'] \approx 0 \to (x \to y)$
(8) $[x' \to (0 \to y)]' \approx (0 \to x) \to (0 \to y)'$
(9) $0 \to (0 \to x)' \approx 0 \to x'$
(10) $0 \to (x' \to y)' \approx x \to (0 \to y')$
(11) $[(x \to 0') \to y]' \approx (0 \to x) \to y'$

$$\begin{aligned} (12) \ 0 \rightarrow [(0 \rightarrow x) \rightarrow y'] \approx x \rightarrow (0 \rightarrow y') \\ (13) \ 0 \rightarrow (x \rightarrow y) \approx x \rightarrow (0 \rightarrow y) \\ (14) \ (x \rightarrow y) \rightarrow y' \approx y \rightarrow (x \rightarrow y)' \\ (15) \ (x' \rightarrow y) \rightarrow [(0 \rightarrow z) \rightarrow x'] \approx (0 \rightarrow y) \rightarrow [(0 \rightarrow z) \rightarrow x'] \\ (16) \ 0 \rightarrow (x \rightarrow y')' \approx 0 \rightarrow (x' \rightarrow y) \\ (17) \ x \rightarrow (y \rightarrow x') \approx y \rightarrow x' \\ (18) \ [(0 \rightarrow x) \rightarrow y] \rightarrow x \approx y \rightarrow x \\ (19) \ [0 \rightarrow (x \rightarrow y)] \rightarrow x \approx (0 \rightarrow y) \rightarrow x \\ (20) \ (0 \rightarrow x) \rightarrow (0 \rightarrow y) \approx x \rightarrow (0 \rightarrow y) \\ (21) \ x \rightarrow y \approx x \rightarrow (x \rightarrow y) \\ (22) \ [\{x \rightarrow (0 \rightarrow y)\} \rightarrow z]' \approx z \rightarrow [(x \rightarrow y) \rightarrow z]' \\ (23) \ [0 \rightarrow (x \rightarrow y)] \rightarrow y' \approx y \rightarrow (x \rightarrow y)' \\ (24) \ x \rightarrow [(y \rightarrow z) \rightarrow x]' \approx (0 \rightarrow y) \rightarrow [x \rightarrow (z \rightarrow x)'] \\ (25) \ 0 \rightarrow [(0 \rightarrow x) \rightarrow y] \approx x \rightarrow (0 \rightarrow y) \\ (26) \ x \rightarrow (y \rightarrow x)' \approx (y \rightarrow 0') \rightarrow x' \\ (27) \ [(x' \rightarrow y) \rightarrow (z \rightarrow x)'] \rightarrow [(y \rightarrow z) \rightarrow x] \approx (y \rightarrow z) \rightarrow x \\ (28) \ [\{0 \rightarrow (x \rightarrow y)'\} \rightarrow (0 \rightarrow y')']' \approx 0 \rightarrow (x \rightarrow y)' \\ (29) \ [[0 \rightarrow \{(x \rightarrow y) \rightarrow z\}] \rightarrow \{0 \rightarrow (y \rightarrow z)\}']' \approx 0 \rightarrow \{(x \rightarrow y) \rightarrow z\} \\ (30) \ [x \rightarrow (0 \rightarrow y)']' \approx x' \rightarrow (y \rightarrow 0')' \\ (31) \ [(0 \rightarrow x) \rightarrow y]' \approx y \rightarrow (x \rightarrow y)' \\ (32) \ [x \rightarrow (y \rightarrow 0')']' \approx x' \rightarrow (0 \rightarrow y)' \\ (33) \ (x \rightarrow y)' \rightarrow (0 \rightarrow y)' \approx y' \rightarrow x' \\ (34) \ (0 \rightarrow x)' \rightarrow (0 \rightarrow y)' \approx (0 \rightarrow (x' \rightarrow y')) \\ (35) \ [(x \rightarrow y)' \rightarrow \{y \rightarrow (x \rightarrow y)'\}]' \approx (0 \rightarrow x) \rightarrow y \\ (37) \ [\{x \rightarrow (y \rightarrow x)\} \rightarrow (x \rightarrow y)']' \approx (0 \rightarrow x) \rightarrow y \\ (37) \ [\{x \rightarrow (y \rightarrow x)\} \rightarrow (x \rightarrow y)']' \approx (x \rightarrow y)'. \end{aligned}$$

3 Partial order in Implication Zroupoids

Let $\mathbf{A} = \langle A; \rightarrow, 0 \rangle \in \mathbf{I}$. We define the operations \wedge and \vee on \mathbf{A} by:

- $x \wedge y := (x \rightarrow y')'$,
- $x \lor y := (x' \land y')'$.

Note that the above definition of \wedge is a simultaneous generalization of the \wedge operation of De Morgan algebras and that of **SL** (= semilattices with least element 0). It is, of course, well known that the meet operation induces a partial order on both **DM** and **SL**, which naturally leads us to the following definition of a binary relation \sqsubseteq on algebras in **I**.

Definition 3.1 Let $\mathbf{A} \in \mathbf{I}$. We define the relation \sqsubseteq on A as follows:

 $x \sqsubseteq y$ if and only if $x \land y = x$ (equivalently, $(x \to y')' = x$).

For $a, b \in A$, we write

- $a \sqsubset b$ if $a \sqsubseteq b$ and $a \neq b$,
- $a \sqsupseteq b$ if $b \sqsubseteq a$, and
- $a \sqsupset b$ if $a \sqsupset b$ and $a \neq b$.

We already know from [3] that $\langle A; \wedge, \vee \rangle$ is a lattice if and only if **A** is a De Morgan Algebra, implying that \sqsubseteq is a partial order on A. We know (see [3]) that \sqsubseteq is also a partial order on algebras in **SL**. This fact led us naturally to consider the possibility of the existence of a subvariety **V** of **I**, containing both **SL** and **DM**, such that, for every algebra **A** in **V**, the relation \sqsubseteq on **A** is actually a partial order.

In this section we will prove our first main result which says that the subvariety $I_{2,0}$, is a maximal subvariety of I with respect to the property that the relation \sqsubseteq is a partial order on every member of that variety. To achieve this end, we need to, first, prove that \sqsubseteq is indeed a partial order on every member of $I_{2,0}$, which will be done using the following lemmas.

Lemma 3.2 Let $\mathbf{A} \in \mathbf{I}_{2,0}$. Then the relation \sqsubseteq is antisymmetric on \mathbf{A} .

Proof Let $a, b \in A$ such that $a \sqsubseteq b$ and $b \sqsubseteq a$. Let $c \in A$ be arbitrary. Then, using (I) and the hypothesis, one observes that $(c \to a) \to b' = [(b \to c) \to (a \to b')']' = [(b \to c) \to a]'$. Consequently,

(3.1) $(c \to a) \to b' = [(b \to c) \to a]'$, where $c \in A$.

Hence,

$$a' = (a \land b)' \qquad \text{by hypothesis} \\ = (a \to b')'' \qquad \text{by definition of } \land \\ = a \to b' \\ = (a' \to a) \to b' \qquad \text{using Lemma 2.4}(d) \\ = [(b \to a') \to a]' \qquad \text{from (3.1) with } c = a' \\ = [(b \to a')'' \to a]' \\ = (b' \to a)' \qquad \text{by hypothesis,} \end{cases}$$

and, therefore,

(3.2)
$$a' = (b' \to a)'.$$

Now,

$$b' = [b \to a']'' \qquad \text{by hypothesis} \\ = b \to a' \\ = (0 \to a'') \to (b \to a') \qquad \text{by Lemma 2.6 (2) with } x = a', y = b \\ = (0 \to a) \to (b \to a')'' \\ = (0 \to a) \to b' \qquad \text{by hypothesis.}$$

Thus,

$$(3.3) \quad b' = (0 \to a) \to b'.$$

Therefore,

$$a' = [b' \rightarrow a]' \quad \text{from } (3.2)$$

= $[(b \rightarrow 0) \rightarrow a]'$
= $(0 \rightarrow a) \rightarrow b' \quad \text{from } (3.1) \text{ with } c = 0$
= $b' \quad \text{by } (3.3).$

Consequently, we have that a = a'' = b'' = b, thus proving that \sqsubseteq is antisymmetric on **A**. \Box

Now, we turn to proving the transitivity of the relation \sqsubseteq . For this, we need the following lemmas. The proof of the following (technical) lemma is given in the Appendix.

Lemma 3.3 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ with $a, b \in A$ such that $a \sqsubseteq b$. Let $d \in A$ be arbitrary. Then

(1) $(0 \rightarrow a') \rightarrow b = a' \rightarrow b$ (2) $b \rightarrow a' = (0 \rightarrow b) \rightarrow a'$ (3) $b \rightarrow a' = a'$ (4) $0 \rightarrow (a' \rightarrow b) = 0 \rightarrow a$ (5) $[(b \rightarrow d) \rightarrow a]' = (d \rightarrow a) \rightarrow b'$ (6) $(0 \rightarrow d) \rightarrow a' = [\{d \rightarrow (0 \rightarrow b')\} \rightarrow a]'$ (7) $a \rightarrow [(a' \rightarrow d) \rightarrow \{(0 \rightarrow a) \rightarrow b'\}] = (0 \rightarrow d) \rightarrow a'$ (8) $a \rightarrow [(d \rightarrow a) \rightarrow b'] = a \rightarrow (d \rightarrow a)'$ (9) $[0 \rightarrow (b \rightarrow d)] \rightarrow a = (0 \rightarrow d) \rightarrow a$ (10) $[b \rightarrow (a \rightarrow d)] \rightarrow a = (0 \rightarrow d) \rightarrow a$ (11) $b \rightarrow (0 \rightarrow a') = 0 \rightarrow a'$ (12) $[(d \rightarrow a) \rightarrow b']' = (b \rightarrow d) \rightarrow a$ (13) $a' \rightarrow b = b' \rightarrow a$ (14) $(d \rightarrow a') \rightarrow b = (d \rightarrow 0') \rightarrow (a' \rightarrow b)$

- (15) $[(0 \rightarrow a') \rightarrow b]' = (0 \rightarrow a) \rightarrow b'$ (16) $(a' \rightarrow b)' = (0 \rightarrow a) \rightarrow b'$
- (17) $b' \to [(b \to d) \to a] \sqsubseteq 0 \to b.$

Lemma 3.4 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b, e \in A$ such that $(a \to b')' = a$ and $(0 \to e') \to b = b$, and let $d \in A$ be arbitrary. Then

 $(a) \quad b \to d = (0 \to (d \to e)) \to (b \to d)$ $(b) \quad (0 \to e) \to a' = a'$ $(c) \quad (0 \to e') \to a = a.$ $(d) \quad (0 \to e) \to [a \to (a \to d)] = a \to d.$

Proof

(b) Using Lemma 3.3 (3) (twice), and (a) with d = a', we obtain $[0 \to (a' \to e)] \to a' = [0 \to (a' \to e)] \to (b \to a') = b \to a' = a'$. Hence,

(3.4) $[0 \to (a' \to e)] \to a' = a'.$

Then,

$$(0 \rightarrow e) \rightarrow a' = [0 \rightarrow (a' \rightarrow e)] \rightarrow a'$$
 by Lemma 2.6 (19) using $x = a', y = e$
= a' by (3.4).

(c)

$$(0 \to e') \to a = [0 \to (0 \to e)'] \to a \text{ by Lemma 2.6 (9)}$$
$$= [(0 \to e) \to 0'] \to a \text{ by Lemma 2.5 (a)}$$
$$= [(0 \to e) \to a'] \to a \text{ by Lemma 2.6 (1)}$$
$$= a' \to a \text{ by (b)}$$
$$= a \text{ by Lemma 2.4 (d).}$$

z = a

with

Thus,

$$(3.5) \ a \to d = [0 \to (d \to e)] \to (a \to d).$$

Now,

$$\begin{array}{l} (0 \rightarrow e) \rightarrow [a \rightarrow (a \rightarrow d)] = [0 \rightarrow [\{a \rightarrow (a \rightarrow d)\} \rightarrow e]] \rightarrow [a \rightarrow (a \rightarrow d)] \\ & \text{by Lemma 2.6 (19) with } x = a \rightarrow (a \rightarrow d), y = e \\ = [0 \rightarrow [\{a \rightarrow (a \rightarrow d)\} \rightarrow e]] \rightarrow [a \rightarrow \{a \rightarrow (a \rightarrow d)\}] \\ & \text{by Lemma 2.6 (21)} \\ = a \rightarrow [a \rightarrow (a \rightarrow d)] \\ & \text{by (3.5) replacing } d \text{ with } a \rightarrow (a \rightarrow d) \\ = a \rightarrow d \\ & \text{by Lemma 2.6 (21).} \end{array}$$

Thus, (d) is proved and the proof of the lemma is complete.

Each of the next three lemmas prove a crucial step in the proof of transitivity of \sqsubseteq .

Lemma 3.5 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b, c \in A$ such that $a \sqsubseteq b$ and $b \sqsubseteq c$. Let $d, e, f \in A$ be arbitrary. Then

(1)
$$(0 \rightarrow c') \rightarrow b = b$$

(2) $(0 \rightarrow c) \rightarrow [a \rightarrow (a \rightarrow d)] = a \rightarrow d$
(3) $(0 \rightarrow c) \rightarrow (a \rightarrow d) = a \rightarrow d$
(4) $[0 \rightarrow ((0 \rightarrow b) \rightarrow c')] \rightarrow b = b$
(5) $\{d' \rightarrow [0 \rightarrow ((0 \rightarrow b) \rightarrow c')]\} \rightarrow (b \rightarrow d)' = (b \rightarrow d)'$
(6) $(b \rightarrow d) \rightarrow [e \rightarrow (b \rightarrow d)]' = [e \rightarrow 0'] \rightarrow (b \rightarrow d)'$
(7) $[b \rightarrow (a \rightarrow c')] \rightarrow a = a$
(9) $(0 \rightarrow b) \rightarrow (a \rightarrow b) \rightarrow b \rightarrow b$

(8) $(0 \to b) \to (a \to d) = a \to d$

- (9) $0 \to [b \to (a \to d)] = 0 \to (a \to d)$ (10) $0 \to [\{b \to (a \to d)\} \to e] = 0 \to [(a \to d) \to e]$ (11) $[0 \to (d' \to c)] \to (0 \to b)' = (0 \to d) \to (0 \to b)'$ (12) $0 \to (a' \to c) \equiv 0 \to b$ (13) $(0 \to a) \to (0 \to b)' = (0 \to a)'$ (14) $0 \to (a' \to c) = 0 \to a$ (15) $(d \to e) \to [\{b \to (a \to f)\}' \to (0 \to a)'] = (d \to e) \to [(a' \to b) \to (f' \to a')].$ **Proof** By hypothesis, we have $(a \to b')' = a$ and $(b \to c')' = b$.
 - (1)

$$\begin{array}{rcl} (0 \rightarrow c') \rightarrow b &=& (c \rightarrow 0') \rightarrow b & \text{by Lemma 2.5 (a)} \\ &=& [(b' \rightarrow c) \rightarrow (0' \rightarrow b)']' & \text{by (I)} \\ &=& [(b' \rightarrow c) \rightarrow b']' & \text{by Lemma 2.4 (a)} \\ &=& [(0 \rightarrow c) \rightarrow b']' & \text{by Lemma 2.6 (5)} \\ &=& [(c' \rightarrow 0') \rightarrow b']' & \text{by Lemma 2.5 (a)} \\ &=& [(b'' \rightarrow c') \rightarrow (0' \rightarrow b')']'' & \text{from (I)} \\ &=& (b'' \rightarrow c') \rightarrow (0' \rightarrow b')' \\ &=& (b \rightarrow c') \rightarrow b'' & \text{by Lemma 2.4 (a)} \\ &=& (b \rightarrow c') \rightarrow b \\ &=& b' \rightarrow b & \text{by hypothesis} \\ &=& b & \text{by Lemma 2.4 (d).} \end{array}$$

- (2) This is immediate from (1) and Lemma 3.4 (d) with e = c.
- (3) Using Lemma 2.6 (21) and (2) we have that $(0 \to c) \to (a \to d) = (0 \to c) \to [a \to (a \to d)] = a \to d$, implying (3).

$$\begin{bmatrix} 0 \to ((0 \to b) \to c') \end{bmatrix} \to b = \left\{ (b' \to 0) \to [((0 \to b) \to c') \to b]' \right\}' \quad \text{by (I)} \\ = \left\{ b \to [((0 \to b) \to c') \to b]' \right\}' \\ = \left\{ b \to (c' \to b)' \right\}' \qquad \text{by Lemma 2.6 (18)} \\ = \left\{ (b' \to b) \to (c' \to b)' \right\}' \qquad \text{by Lemma 2.4 (d)} \\ = (b \to c') \to b \qquad \text{by (I)} \\ = (b \to c')'' \to b \\ = b' \to b \qquad \text{by hypothesis} \\ = b \qquad \text{by Lemma 2.4 (d)}.$$

(5)

$$\{d' \to [0 \to ((0 \to b) \to c')]\} \to (b \to d)' = \{[[0 \to ((0 \to b) \to c')] \to b] \to d\}' \text{ by (I)}$$
$$= (b \to d)' \text{ by (4)}.$$

$$\begin{aligned} (b \to d) \to [e \to (b \to d)]' &= [e \to (b \to d)] \to (b \to d)' & \text{by Lemma 2.6 (14) with} \\ & x = e, y = b \to d \\ &= [e \to 0'] \to (b \to d)' & \text{by Lemma 2.6 (1).} \end{aligned}$$

(8)

$$(0 \rightarrow b') \rightarrow b = (0 \rightarrow 0') \rightarrow b \quad \text{by Lemma 2.6 (1)}$$
$$= (0'' \rightarrow 0') \rightarrow b$$
$$= 0' \rightarrow b \qquad \text{by Lemma 2.4 (d)}$$
$$= b \qquad \text{by Lemma 2.4 (a)}.$$

Hence, by the hypothesis, together with Lemma 3.4 (d), we obtain that $(0 \to b) \to \{a \to (a \to d)\} = a \to d$. Hence, by Lemma 2.6 (21), we have $(0 \to b) \to (a \to d) = a \to d$.

$$\begin{array}{rcl} 0 \to (a \to d) &=& 0 \to \left[(0 \to b) \to (a \to d) \right] & \text{by (8)} \\ &=& b \to \left[0 \to (a \to d) \right] & \text{by Lemma 2.6 (25) with } x = b, y = a \to d \\ &=& 0 \to \left[b \to (a \to d) \right]. & \text{by Lemma 2.6 (13).} \end{array}$$

$$\begin{array}{rcl} 0 \rightarrow [\{b \rightarrow (a \rightarrow d)\} \rightarrow e] &=& [b \rightarrow (a \rightarrow d)] \rightarrow (0 \rightarrow e) & \text{by Lemma 2.6 (13)} \\ &=& 0 \rightarrow [[0 \rightarrow \{b \rightarrow (a \rightarrow d)\}] \rightarrow e] & \text{by Lemma 2.6 (25)} \\ &=& 0 \rightarrow [\{0 \rightarrow (a \rightarrow d)\} \rightarrow e] & \text{by (9)} \\ &=& (a \rightarrow d) \rightarrow (0 \rightarrow e) & \text{by Lemma 2.6 (25)} \\ &=& 0 \rightarrow [(a \rightarrow d) \rightarrow e] & \text{by Lemma 2.6 (13)}. \end{array}$$

(11)

(12)

$$\begin{array}{rcl}
0 \to (a' \to c) &=& 0 \to [(a \to b')'' \to c] & \text{by hyphotesis} \\
 &=& 0 \to [(a \to b') \to c] \\
 &\sqsubseteq& 0 \to (b' \to c) & \text{by Lemma 2.6 (29)} \\
 &=& 0 \to b & \text{by hyphotesis and Lemma 3.3 (4).} \end{array}$$

$$(0 \to a) \to (0 \to b)' = [a' \to (0 \to b)]' \text{ by Lemma 2.6 (8)}$$

= $[0 \to (a' \to b)]'$ by Lemma 2.6 (13)
= $(0 \to a)'$. by hyphotesis and Lemma 3.3 (4).

(14)

$$\begin{array}{rcl} 0 \rightarrow (a' \rightarrow c) &=& [\{0 \rightarrow (a' \rightarrow c)\} \rightarrow (0 \rightarrow b)']' & \text{by (12)} \\ &=& [(0 \rightarrow a) \rightarrow (0 \rightarrow b)']' & \text{by (11) with } d = a \\ &=& (0 \rightarrow a)'' & \text{by (13)} \\ &=& 0 \rightarrow a. \end{array}$$

(15)

Hence, we have $(d \to e) \to [\{b \to (a \to f)\}' \to (0 \to a)'] = (d \to e) \to [(a' \to b) \to (f' \to a')].$

Lemma 3.6 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b, c \in A$ such that $a \sqsubseteq b$ and $b \sqsubseteq c$. Let $d \in A$ be arbitrary. Then

(a) $[c \rightarrow (b \rightarrow a')] \rightarrow b = (0 \rightarrow a') \rightarrow b$ (b) $(c \rightarrow a') \rightarrow b = a' \rightarrow b$ (c) $(a' \rightarrow b) \rightarrow (c \rightarrow a') = c \rightarrow a'$ (d) $c \rightarrow a' = a \rightarrow [b \rightarrow (a \rightarrow c')]$ (e) $0 \rightarrow (a \rightarrow d) = 0 \rightarrow [c \rightarrow (a \rightarrow d)]$

(f)
$$(d \to a) \to d \sqsubseteq (a' \to b) \to d$$

(g) $(a' \to b) \to c' = (0 \to a) \to b'$
(h) $0 \to (a \to c') \sqsubseteq 0 \to a'$
(i) $0 \to (a \to c') \equiv 0 \to a'$.
(j) $c \to (a \to c') \sqsubseteq 0 \to (a \to c')$
(k) $c \to (a \to c') \sqsubseteq 0 \to a'$
(l) $(c \to (a \to c'))' \to (0 \to a)' = c \to (a \to c')$
(m) $a \to [b \to (a \to c')] = a \to c'$
(n) $c \to a' = a \to c'$.

Proof

- (a) Since $(b \to c')' = b$, by Lemma 3.3 (10) with d = a', we have $(c \to (b \to a')) \to b = (0 \to a') \to b$.
- (b)

$$(c \to a') \to b = [c \to (b \to a')] \to b \text{ by Lemma 3.3 (3)} = (0 \to a') \to b \text{ by (a)},$$

from which we get $(c \to a') \to b = (0 \to a') \to b$, which, together with Lemma 3.3 (1), implies $(c \to a') \to b = a' \to b$.

(c)

$$c \rightarrow a' = (0 \rightarrow a) \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (2) with $x = a', y = c$

$$= [0 \rightarrow (a' \rightarrow b)] \rightarrow (c \rightarrow a')$$
by Lemma 3.3 (4)

$$= [0 \rightarrow \{(c \rightarrow a') \rightarrow b\}] \rightarrow (c \rightarrow a')$$
by (b)

$$= [(c \rightarrow a') \rightarrow (0 \rightarrow b)] \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (13)

$$= [0 \rightarrow (0 \rightarrow b)] \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (5)

$$= (0 \rightarrow b) \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (6)

$$= [(c \rightarrow a') \rightarrow b] \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (5)

$$= (a' \rightarrow b) \rightarrow (c \rightarrow a')$$
by Lemma 2.6 (5)

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$$\begin{array}{lll} c \rightarrow a' &=& (0 \rightarrow a) \rightarrow (c \rightarrow a') & \text{by Lemma 2.6 (2)} \\ &=& (0 \rightarrow a) \rightarrow [(a' \rightarrow b) \rightarrow (c'' \rightarrow a')] & \text{by } (c) \\ &=& (0 \rightarrow a) \rightarrow [(b \rightarrow (a \rightarrow c'))' \rightarrow (0 \rightarrow a)'] & \text{by Lemma 3.5 (15) with} \\ &=& (0 \rightarrow a) \rightarrow [\{b \rightarrow (a \rightarrow c')\}' \rightarrow \{0 \rightarrow (a' \rightarrow c)\}'] & \text{by Lemma 3.5 (15) with} \\ &=& (0 \rightarrow a) \rightarrow [\{b \rightarrow (a \rightarrow c')\}' \rightarrow \{0 \rightarrow (a \rightarrow c')\}'] & \text{by Lemma 3.5 (14)} \\ &=& (0 \rightarrow a) \rightarrow [\{b \rightarrow (a \rightarrow c')\}' \rightarrow [0 \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 3.5 (10) with} \\ &=& (0 \rightarrow a) \rightarrow [\{b \rightarrow (a \rightarrow c')\}' \rightarrow [0 \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 2.6 (16)} \\ &=& (0 \rightarrow a) \rightarrow [\{b \rightarrow (a \rightarrow c')\}' \rightarrow [0 \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 2.6 (24) with} \\ &=& [b \rightarrow (a \rightarrow c')]' \rightarrow [(a \rightarrow 0) \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 2.6 (24) with} \\ &=& [b \rightarrow (a \rightarrow 0)] \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 2.6 (4) and (5)} \\ &=& [\{a \rightarrow (0 \rightarrow 0)\} \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 2.6 (13)} \\ &=& [\{a \rightarrow (0 \rightarrow 0)\} \rightarrow \{b \rightarrow (a \rightarrow c')\}']' & \text{by Lemma 3.5(6) with} \\ &=& a, d = a \rightarrow c' \\ &=& a \rightarrow [b \rightarrow (a \rightarrow c')] & \text{by Lemma 3.5(7)}. \end{array}$$

$$\begin{array}{rcl} 0 \rightarrow (a \rightarrow d) &=& 0 \rightarrow \left[(0 \rightarrow c) \rightarrow (a \rightarrow d) \right] & \text{by Lemma 3.5(3)} \\ &=& c \rightarrow \left[0 \rightarrow (a \rightarrow d) \right] & \text{by Lemma 2.6 (25)} \\ &=& 0 \rightarrow \left[c \rightarrow (a \rightarrow d) \right] & \text{by Lemma 2.6 (13).} \end{array}$$

(f)

(g)

$$\begin{array}{rcl} (d \rightarrow a) \rightarrow d &=& (0 \rightarrow a) \rightarrow d & \qquad \text{by Lemma 2.6 (5)} \\ &=& [0 \rightarrow (a' \rightarrow b)] \rightarrow d & \qquad \text{by Lemma 3.3 (4)} \\ &\sqsubseteq& (a' \rightarrow b) \rightarrow d & \qquad \text{by Lemma 2.6 (36).} \end{array}$$

$$\begin{aligned} (a' \to b) \to c' &= [(c \to a') \to (b \to c')']' & \text{by (I)} \\ &= [(c \to a') \to b]' & \text{by hypothesis} \\ &= [\{c \to (b \to a')\} \to b]' & \text{by Lemma 3.3 (3)} \\ &= [(0 \to a') \to b]' & \text{by Lemma 3.3 (10) with} \\ &\quad d = a' \text{ since } b \sqsubseteq c \\ &= [(b \to a') \to b]' & \text{by Lemma 2.6 (5)} \\ &= [(b \to a') \to b'']' \\ &= [(b \to a') \to (0' \to b')']' & \text{by Lemma 2.4 (a)} \\ &= (a' \to 0') \to b' & \text{by (I)} \\ &= (0 \to a) \to b'. & \text{by Lemma 2.5 (a).} \end{aligned}$$

Hence, one has $(a' \to b) \to c' = (0 \to a) \to b'$.

(h) From Lemma 3.5 (1), we have $(0 \to c') \to b = b$. Hence, we can use Lemma 3.3. Therefore, we have

$$\begin{array}{rcl} 0 \rightarrow (a \rightarrow c') &=& 0 \rightarrow [\{(0 \rightarrow c') \rightarrow a\} \rightarrow c'] & \text{ by Lemma 3.4 (c) and Lemma 3.5(1)} \\ &=& [(0 \rightarrow c') \rightarrow a] \rightarrow (0 \rightarrow c') & \text{ by Lemma 2.6 (13)} \\ &\sqsubseteq& (a' \rightarrow b) \rightarrow (0 \rightarrow c') & \text{ by (f) with } d = 0 \rightarrow c' \\ &=& 0 \rightarrow [(a' \rightarrow b) \rightarrow c'] & \text{ by Lemma 2.6 (13)} \\ &=& 0 \rightarrow [(0 \rightarrow a) \rightarrow b'] & \text{ by (g)} \\ &=& 0 \rightarrow [(b \rightarrow 0) \rightarrow (a \rightarrow b')']' & \text{ by (I)} \\ &=& 0 \rightarrow [b' \rightarrow (a \rightarrow b')']' & \text{ by (I)} \\ &=& 0 \rightarrow (b' \rightarrow a)' & \text{ by hypothesis} \\ &=& 0 \rightarrow (b \rightarrow a') & \text{ by Lemma 2.6 (16)} \\ &=& 0 \rightarrow a' & \text{ by Lemma 3.3 (3).} \end{array}$$

$$\begin{array}{rcl} 0 \to a' &=& 0 \to (a \to 0) \\ &=& 0 \to [c \to (a \to 0)] & \text{by (e)} \\ &=& 0 \to (c \to a') \\ &=& 0 \to [(a \to c')' \to (0 \to a)'] & \text{by Lemma 2.6 (33)} \\ &=& [0 \to (a \to c')]' \to (0 \to a)' & \text{by Lemma 2.6 (34) and Lemma 2.6 (6)} \\ &=& [0 \to (a \to c')]' \to (a' \to 0')' & \text{by Lemma 2.5 (a)} \\ &=& [\{0 \to (a \to c')\} \to (0 \to a')']' & \text{by Lemma 2.5 (30) with } x = 0 \to (a \to c'), y = a' \\ &=& 0 \to (a \to c') & \text{by (h).} \end{array}$$

$$\begin{aligned} \text{(j)} \\ & [\{c \rightarrow (a \rightarrow c')\} \rightarrow \{0 \rightarrow (a \rightarrow c')\}']' = [\{0 \rightarrow (a \rightarrow c')\} \rightarrow c] \rightarrow [(a \rightarrow c') \rightarrow \{0 \rightarrow (a \rightarrow c')\}']' \\ & \text{by (I)} \\ & = [\{0 \rightarrow (a \rightarrow c')\} \rightarrow c] \rightarrow [\{(a \rightarrow c') \rightarrow 0\} \rightarrow (a \rightarrow c')] \\ & \text{by Lemma 2.6 (4)} \\ & = [\{0 \rightarrow (a \rightarrow c')\} \rightarrow c] \rightarrow [(a \rightarrow c')' \rightarrow (a \rightarrow c')] \\ & = [\{0 \rightarrow (a \rightarrow c')\} \rightarrow c] \rightarrow (a \rightarrow c') \\ & \text{by Lemma 2.4 (d)} \\ & = c \rightarrow (a \rightarrow c') \\ \end{aligned}$$

(k) From (j) we have that $c \to (a \to c') \sqsubseteq 0 \to (a \to c')$. Then using (i) we get $c \to (a \to c') \sqsubseteq 0 \to a'$.

$$[c \to (a \to c')]' \to (0 \to a)' = [c \to (a \to c')]' \to (a' \to 0')' \text{ by Lemma 2.5 (a)}$$
$$= [\{c \to (a \to c')\} \to (0 \to a')']' \text{ by Lemma 2.6 (30)}$$
$$= c \to (a \to c') \text{ by (k).}$$

(m)

Lemma 3.7 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b, c \in A$ such that $a \sqsubseteq b$ and $b \sqsubseteq c$. Then

(a)
$$c' \rightarrow [(c \rightarrow d) \rightarrow b] \sqsubseteq c$$

(b) $0 \rightarrow a' = c \rightarrow (0 \rightarrow a')$
(c) $c' \rightarrow (a' \rightarrow b) \sqsubseteq c$
(d) $(0 \rightarrow a') \rightarrow b = (c \rightarrow a') \rightarrow b$
(e) $c' \rightarrow (a' \rightarrow b) \sqsubseteq 0 \rightarrow c$
(f) $[(0 \rightarrow a) \rightarrow b'] \rightarrow c = c' \rightarrow (a' \rightarrow b)$
(g) $a' \rightarrow c = c' \rightarrow (a' \rightarrow b)$
(h) $a' \rightarrow c \sqsubseteq c$
(i) $a' \rightarrow c = (0 \rightarrow a') \rightarrow c$.

Proof

(a)

$$c' \rightarrow [(c \rightarrow d) \rightarrow b] = c' \rightarrow [(c \rightarrow d) \rightarrow (b \rightarrow c')'] \text{ by hypothesis}$$
$$= c' \rightarrow [(d \rightarrow b) \rightarrow c']' \text{ by (I)}$$
$$\sqsubseteq c'' \text{ by Lemma 2.6 (37)}$$
$$= c.$$

$$\begin{array}{rcl} 0 \rightarrow a' &=& b \rightarrow (0 \rightarrow a') & \text{by Lemma 3.3 (11)} \\ &=& [0 \rightarrow \{(0 \rightarrow a') \rightarrow c\}] \rightarrow [b \rightarrow (0 \rightarrow a')] & \text{by Lemma 3.5 (1)} \\ && \text{and Lemma 3.4 (a)} \\ && \text{with } d = 0 \rightarrow a', e = c \\ &=& [0 \rightarrow \{(0 \rightarrow a') \rightarrow c\}] \rightarrow (0 \rightarrow a') & \text{by Lemma 3.3 (11)} \\ &=& [(0 \rightarrow a') \rightarrow (0 \rightarrow c)] \rightarrow (0 \rightarrow a') & \text{by Lemma 2.6 (13)} \\ &=& [0 \rightarrow (0 \rightarrow c)] \rightarrow (0 \rightarrow a') & \text{by Lemma 2.6 (5)} \\ &=& (0 \rightarrow c) \rightarrow (0 \rightarrow a') & \text{by Lemma 2.6 (21)} \\ &=& c \rightarrow (0 \rightarrow a') & \text{by Lemma 2.6 (20).} \end{array}$$

(c)

$$\begin{array}{lll} (0 \rightarrow a') \rightarrow b &=& [c \rightarrow (0 \rightarrow a')] \rightarrow b & \text{by (b)} \\ &=& [(b' \rightarrow c) \rightarrow \{(0 \rightarrow a') \rightarrow b\}']' & \text{by (I)} \\ &=& [(b' \rightarrow c) \rightarrow \{(b' \rightarrow 0) \rightarrow (a' \rightarrow b)'\}]' & \text{by (I) and } x'' \approx x \\ &=& [(b' \rightarrow c) \rightarrow \{b \rightarrow (a' \rightarrow b)'\}]' & \text{by Lemma 2.6 (26)} \\ &=& [(b' \rightarrow c) \rightarrow \{(0 \rightarrow a) \rightarrow b'\}]' & \text{by Lemma 2.5 (a)} \\ &=& [(b' \rightarrow c) \rightarrow (a' \rightarrow b)']' & \text{by Lemma 3.3 (16)} \\ &=& (c \rightarrow a') \rightarrow b & \text{by (I)}. \end{array}$$

(e)

(d)

$$\begin{array}{rcl} c' \to (a' \to b) &=& c' \to [(0 \to a') \to b] & \text{by Lemma 3.3 (15) and Lemma 3.3 (16)} \\ &=& c' \to [(c \to a') \to b] & \text{by (d)} \\ &\sqsubseteq& 0 \to c. & \text{by Lemma 3.3 (17) with } d = a'. \end{array}$$

(f)

$$\begin{aligned} c' \to (a' \to b) &= [\{c' \to (a' \to b)\} \to (0 \to c)']' & \text{by (e)} \\ &= [(a' \to b) \to 0] \to c & \text{by (I)} \\ &= (a' \to b)' \to c \\ &= [(0 \to a) \to b'] \to c & \text{by Lemma 3.3 (16).} \end{aligned}$$

(g)

$$\begin{aligned} c' \to (a' \to b) &= ((0 \to a) \to b') \to c & \text{by (f)} \\ &= [(0 \to a) \to 0'] \to (b' \to c) & \text{by Lemma 3.3 (14) with } d = 0 \to a \\ &= [(a' \to 0') \to 0'] \to (b' \to c) & \text{by Lemma 2.5 (a)} \\ &= [(a' \to 0) \to 0'] \to (b' \to c) & \text{by Lemma 2.6 (1)} \\ &= [a'' \to 0'] \to (b' \to c) \\ &= (a \to 0') \to (b' \to c) \\ &= (a \to b') \to c & \text{by Lemma 3.3 (14) with } d = a \\ &= a' \to c & \text{by hypothesis.} \end{aligned}$$

(h) This is immediate from (g) and (c).

(i)

$$(0 \to a') \to c = (c \to a') \to c \quad \text{by Lemma 2.6 (5)} \\ = [c \to (a' \to c)']' \quad \text{by Lemma 2.6 (4)} \\ = [(a' \to c) \to c']' \quad \text{by Lemma 2.6 (14)} \\ = a' \to c \qquad \text{by (h).}$$

Theorem 3.8 \sqsubseteq *is transitive.*

We are now ready to complete the proof of transitivity of $\sqsubseteq.$

Proof Let $a, b, c \in A$ such that $a \sqsubseteq b$ and $b \sqsubseteq c$. Observe that

 $a' = a \rightarrow 0$ $= (0 \rightarrow c) \rightarrow (a \rightarrow 0) \quad \text{by Lemma 3.5 (3) with } d = 0$ $= (0 \rightarrow c) \rightarrow a'$ $= (a' \rightarrow c) \rightarrow a' \quad \text{by Lemma 2.6 (5)}$ $= ((0 \rightarrow a') \rightarrow c) \rightarrow a' \quad \text{by Lemma 3.7 (i)}$ $= c \rightarrow a' \quad \text{by Lemma 2.6 (18)}$ $= a \rightarrow c' \quad \text{by Lemma 3.6 (n).}$

Consequently,

$$a = a'' = (a \to c')',$$

implying $a \sqsubseteq c$. Hence, \sqsubseteq is transitive on **A**.

We are now prepared to present our first main theorem.

Theorem 3.9 The variety $I_{2,0}$ is a maximal subvariety of I with respect to the property that the relation \sqsubseteq introduced in Definition 3.1 is a partial order.

Proof Let $\mathbf{A} \in \mathbf{I}_{2,0}$. The relation \sqsubseteq is a partial order on A in view of Lemma 2.4 (c), Lemma 3.2, and Theorem 3.8.

Next, let **V** be a subvariety of **I** such that \sqsubseteq is a partial order on every algebra in **V**. Now let $\mathbf{A} \in \mathbf{V}$. Reflexivity of \sqsubseteq implies that $\mathbf{A} \models (x \rightarrow x')' \approx x$. Therefore, by Lemma 2.4, we conclude that $\mathbf{A} \in \mathbf{I}_{2,0}$, and hence, $\mathbf{V} \subseteq \mathbf{I}_{2,0}$, completing the proof.

4 A method to construct finite $I_{2,0}$ -chains

Now that we know the relation \sqsubseteq is a partial order on algebras in $\mathbf{I}_{2,0}$, it is natural to consider those algebras in $\mathbf{I}_{2,0}$, in which \sqsubseteq is a total order.

Definition 4.1 Let $\mathbf{A} \in \mathbf{I}$. We say that \mathbf{A} is an $\mathbf{I}_{2,0}$ -chain (chain, for short) if $\mathbf{A} \in \mathbf{I}_{2,0}$ and the relation \sqsubseteq (see Definition 3.1) is totally ordered on A.

In this section we describe a method of constructing finite $I_{2,0}$ -chains. But, first, we will present some examples of $I_{2,0}$ -chains that will foreshadow the method to construct finite $I_{2,0}$ -chains. We note that, in these examples, the number 0 is the constant element.

It is easy to see that the only 2-element $I_{2,0}$ -chains, up to isomorphism, are

and the only 3-element $I_{2,0}$ -chains, up to isomorphism, are

Note that, henceforth, we will use the symbol \leq to denote the natural order in \mathbb{Z} . Recall that \sqsubseteq is being used for the order (see Definition 3.1).

The next definition describes a general method to construct a class of finite $I_{2,0}$ -chains, generalizing the above examples. In the next section, we will show that, this method, in fact, yields, up to isomorphism, all finite $I_{2,0}$ -chains.

Definition 4.2 Let $k \in \mathbb{N}$. Let $m, n \in \omega$ be such that the interval $[-n, m] \subseteq \mathbb{Z}$ with |[-n, m]| = kand $0 \leq n, m \leq k - 1$. The (auxiliary) functions p (predecessor) and * are defined on [-n, m]as follows:

$$p(x) = \begin{cases} x - 1 & \text{if } x > -n \\ -n & \text{if } x = -n, \end{cases}$$

and

$$x^* = \left\{ \begin{array}{ll} m & \mbox{if} \ x = 0 \\ x & \mbox{if} \ x < 0 \\ p((p(x))^*) & \mbox{if} \ x > 0. \end{array} \right.$$

For convenience, we write $p(p(x)^*)$ for $p((p(x))^*)$. (Notice that the function * is defined recursively for $x \ge 0$.)

Define the algebra [-n, m] as follows:

$$[-\mathbf{n},\mathbf{m}] := \langle [-n,m]; \Rightarrow, 0 \rangle$$
, where $0 \in [-n,m]$ is the constant and \Rightarrow is defined by

$$x \Rightarrow y = \begin{cases} max(x^*, y) & \text{if } x, y \ge 0\\ min(x, y) & \text{otherwise.} \end{cases}$$

We set $x' := x \Rightarrow 0$.

We shall now illustrate the method described in the above definition by applying it to construct a 6-element $I_{2,0}$ -chain.

Let k = 6, and consider the interval $A = [-2,3] = \{-2,-1,0,1,2,3\}$. Since $0 \Rightarrow 0 = max(0^*,0) = max(3,0) = 3$ and $a \Rightarrow b = min(a,b)$ if a < 0 or b < 0, we arrive at the following partial table for \Rightarrow :

\Rightarrow	-2	-1	0	1	2	3
-2	-2	-2	-2	-2	-2	-2
-1	-2	-1	-1	-1	-1	-1
0	-2	-1	3	?	?	?
1	-2	-1	?	?	?	?
2	-2	-1	?	?	?	?
3	-2	-2 -1 -1 -1 -1 -1 -1	?	?	?	?

Next, we determine the operations p and *:

$x \mid x^*$	$x \mid x \Rightarrow 0$
$\begin{array}{c cccc} 0 & 3 \\ 1 & p(p(1)^*) = p(0^*) = p(3) = 2 \\ 2 & p(p(2)^*) = p(1^*) = p(2) = 1 \\ \end{array}$	$\begin{array}{c c} x & x \to 0 \\ \hline 1 & max(1^*, 0) = max(2, 0) = 2 \\ 2 & max(2^*, 0) = max(1, 0) = 1 \\ 3 & max(3^*, 0) = max(0, 0) = 0 \end{array}$
$3 \mid p(p(3)^*) = p(2^*) = p(1) = 0$ Hence the table for \Rightarrow becomes: \Rightarrow -2 -1 0 1 2 3	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	

Observe that $0 \Rightarrow 1 = max(0^*, 1) = max(3, 1) = 3$, $1 \Rightarrow 1 = max(1^*, 1) = max(2, 1) = 2$, $2 \Rightarrow 1 = max(2^*, 1) = max(1, 1) = 1$ and $3 \Rightarrow 1 = max(3^*, 1) = max(0, 1) = 1$. Then we get

\Rightarrow	-2	-1	0	1	2	3
-2	-2	-2	-2	-2		-2
-1	-2	-1	-1	-1	-1	-1
0	-2	-1	3	3	?	?
1	-2	-1	2	2	?	?
2	-2	-1	1	1	?	?
3	-2	-1	0	1	?	?

??

?

?

-2

 $\begin{array}{c|cc}
2 & -2 \\
3 & -2
\end{array}$

-1

-1 -1

1

Iterating this process we obtain the following complete table for \Rightarrow :

\Rightarrow	-2	-1	0	1	2	3
-2	-2	-2	-2	-2	-2	-2
-1	-2	-1	-1	-1	-1	-1
0	-2		3	3	3	3
1	-2	-1	2	2	2	3
2	-2	-1	1	1	2	3
3	-2	-1	0	1	2	3

Thus we have constructed the algebra $[-\mathbf{n}, \mathbf{m}]$. Observe that $-2 \sqsubset -1 \sqsubset 0 \sqsubset 1 \sqsubset 2 \sqsubset 3$ and $x'' = x^{**} = x$. Also, it is routine to verify $[-\mathbf{n}, \mathbf{m}] \in \mathbf{I}_{2,0}$. Hence it is an $\mathbf{I}_{2,0}$ -chain.

Returning to the general method, we now aim to prove that $[-\mathbf{n}; \mathbf{m}]$ is an $\mathbf{I}_{2,0}$ -chain. To prove this, we will need the following lemmas.

Lemma 4.3 If $x \in [-n, m]$ and $0 \le x \le m$ then $x^* = m - x$ and, consequently, $x^* \in [0, m]$.

Proof We prove this lemma by induction on the element x. Assume that x = 0. Then $0^* = m = m - 0$.

Next, suppose x > 0. Since $-n \le 0 < x$, we have p(x) = x - 1. Hence, by inductive hypothesis, we have

$$(4.1) \ p(x)^* = m - p(x) = m - (x - 1) = m - x + 1$$

From x > 0, we can conclude that $m - x + 1 \le m$. Also, since $x \le m$, we obtain $0 \le m - x$, thus $-n - 1 < 0 \le m - x$, implying m - x + 1 > -n. Then we get p(m - x + 1) = m - x + 1 - 1. By (4.1), $x^* = p((p(x))^*) = p(m - x + 1) = m - x$, completing the induction. It is clear that $x^* \in [0, m]$.

Corollary 4.4 If $x \in [-n, m]$ then $x' = x^*$.

Proof If x < 0 we have that $x' = x \Rightarrow 0 = min(x, 0) = x = x^*$. If x > 0, then by Lemma 4.3, $x^* \ge 0$, and hence $x' = x \Rightarrow 0 = max(x^*, 0) = x^*$.

Lemma 4.5 If $x \in [-n, m]$ then x'' = x.

Proof We consider the following cases:

- If x < 0, then $x^* = x$, and hence $x^{**} = x$.
- If $x \ge 0$,

$$x^{**} = (m-x)^*$$
 by Lemma 4.3 since $0 < x \le m$
= $m - (m-x)$ by Lemma 4.3 since $0 \le m - x \le m$
= x .

Consequently, by Corollary 4.4, x'' = x.

Lemma 4.6 If $x, y \in [-n, m]$ and $0 \le x \le y$ then $x^* \ge y^*$.

Proof We prove this lemma by induction on the element x. If x = 0, $x^* = 0^* = m \ge y^*$ by Lemma 4.3.

Now assume that x > 0. Since $0 < x \le y$, we have that $x^* = p(p(x)^*)$ and $y^* = p(p(y)^*)$. Note that $0 \le p(x) \le p(y)$. Then, by induction hypothesis, we get $p(y)^* \le p(x)^*$. Hence $x^* = p(p(x)^*) \ge p(p(y)^*) = y^*$.

Lemma 4.7 Let $k \in \mathbb{N}$. Let $m, n \in \omega$ be such that the interval $[-n, m] \subseteq \mathbb{Z}$ with |[-n, m]| = kand $0 \leq n, m \leq k - 1$. Then, $[-n, m] \in \mathbf{I}_{2,0}$.

Proof The proof that $\langle [-n; m]; \Rightarrow, 0 \rangle$ satisfies the identity (I) is long and computational, but routine. Hence we leave the verification to the reader with the recommendation that the following cases be considered, where $i, j, k \in [-n; m]$:

(7) $i \ge 0, j < 0$ and $k \ge 0$
(8) $i \ge 0, j < 0$ and $k < 0$
(9) $i < 0, j \ge 0$ and $k \ge 0$
(10) $i < 0, j \ge 0$ and $k < 0$
(11) $i < 0, j < 0$ and $k \ge 0$
(12) $i, j, k < 0.$

Observe that, if $x \in [-n, m]$, then, from Corollary 4.4, we have $x' = x^*$, and from Lemma 4.5 we have that x'' = x; and in particular 0'' = 0. Thus, we conclude that $\langle [-n, m]; \Rightarrow, 0 \rangle \in \mathbf{I}_{2,0}$. \Box

In view of the above lemma and Theorem 3.8, the relation defined by

$$x \sqsubseteq y$$
 if and only if $(x \Rightarrow y')' = x$

is a partial order on [-n, m]. We now wish to show that \sqsubseteq is indeed a total order.

Lemma 4.8 Let [-n, m] be the algebra, as defined in Definition 4.2. Then

$$\langle [-n,m]; \sqsubseteq \rangle \cong \langle [-n,m]; \le \rangle.$$

Proof Let $x, y \in [-n, m]$. It is enough to prove that $x \leq y$ if and only if $x \sqsubseteq y$. Assume that $x \leq y$. We will consider the following cases:

• Case 1: x < 0. Then

(4.2)
$$(x \Rightarrow y')' = (x \Rightarrow y^*)^* = [min(x, y^*)]^*.$$

We consider further the following subcases:

- Case 1.1: y < 0.

$$(x \Rightarrow y')' = [min(x, y^*)]^* \quad \text{by (4.2)}$$

= $[min(x, y)]^* \quad \text{since } y < 0$
= $x^* \qquad \text{since } x \le y$
= x . since $x < 0$

- Case 1.2: $y \ge 0$.

$$\begin{array}{rcl} (x \Rightarrow y')' &=& [min(x,y^*)]^* & \mbox{by (4.2)} \\ &=& x^* & \mbox{since } y^* \geq 0 \mbox{ by Lemma 4.3, and } x < 0 \\ &=& x. \end{array}$$

• Case 2: $x \ge 0$. Therefore $y \ge 0$. In this case

$$(x \Rightarrow y')' = (x \Rightarrow y^*)^*$$

= $[max(x^*, y^*)]^*$
= x^{**} by Lemma 4.6
= x

Thus, in all these cases, $x \sqsubseteq y$.

For the converse, suppose $x \sqsubseteq y$.

• Case 1: x < 0. If $y \ge 0$ then x < y. So, we can assume y < 0. Then

$$\begin{array}{rcl} x &=& x' & \text{since } x < 0 \\ &=& (x \Rightarrow y')'' & \text{by hypothesis} \\ &=& x \Rightarrow y' & \text{by Lemma 4.5} \\ &=& x \Rightarrow y \\ &=& \min(x, y). \end{array}$$

Hence $x \leq y$.

• Case 2: $x \ge 0$. Suppose y < 0. Then

$$\begin{aligned} x &= (x \Rightarrow y')' & \text{by hypothesis} \\ &= (x \Rightarrow y)' \\ &= \min(x, y)' \\ &= y' \\ &= y, \end{aligned}$$

a contradiction. Hence $y \ge 0$. Consequently,

$$\begin{aligned} x' &= (x \Rightarrow y')'' \\ &= x \Rightarrow y' & \text{by Lemma 4.5} \\ &= max(x', y'), \end{aligned}$$

so, $x' \ge y'$. Then, by Lemma 4.5 and Lemma 4.6, $x = x'' \le y'' = y$.

In view of Lemma 4.7 and Lemma 4.8, we have proved the following

Theorem 4.9 [-n, m] is an $I_{2,0}$ -chain, where

$$-n \sqsubset -n + 1 \sqsubset \ldots \sqsubset -1 \sqsubset 0 \sqsubset 1 \sqsubset 2 \sqsubset \ldots \sqsubset m.$$

5 Characterization of finite $I_{2,0}$ -chains

In this section we are going to prove our second main result. The following lemmas will be useful later in this section.

Lemma 5.1 Let $\mathbf{A} \in \mathbf{I}_{2,0}$. Then 0' is the greatest element in A, relative to \sqsubseteq .

Proof Let $a \in A$. Since $(a \to (0 \to 0)')' = (a \to 0'')' = (a \to 0)' = a'' = a$, we have $a \sqsubseteq 0'$. \Box

Lemma 5.2 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b \in A$ with $0 \sqsubseteq a \sqsubseteq b$. Then $b' \sqsubseteq a'$.

Proof

$$\begin{aligned} (b' \to a'')' &= (b' \to a)' \\ &= (b' \to (a \to b')')' & \text{by hypothesis} \\ &= ((a \to 0') \to b'')' & \text{by Lemma 2.6 (26)} \\ &= ((a \to 0') \to b)' \\ &= ((0 \to a') \to b)' & \text{by Lemma 2.5 (a)} \\ &= ((0 \to a')'' \to b)' \\ &= (0' \to b)' & \text{by hypothesis} \\ &= b' & \text{by Lemma 2.4 (a).} \end{aligned}$$

Lemma 5.3 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a \in A$. If $0 \sqsubseteq a$ then $0 \rightarrow a = 0'$.

Proof First notice that, since $0 \sqsubseteq a$, $0' = (0 \rightarrow a')'' = 0 \rightarrow a'$. Consequently,

(5.1) $0' = 0 \to a'.$

Then

$$\begin{array}{rcl} 0' &=& 0' \to 0' & \text{by Lemma 2.4 (a)} \\ &=& (0 \to a') \to 0' & \text{by (5.1)} \\ &=& (0' \to a') \to 0' & \text{by Lemma 2.6 (5)} \\ &=& a' \to 0' & \text{by Lemma 2.4 (a)} \\ &=& 0 \to a. & \text{by Lemma 2.5 (a)} \end{array}$$

Lemma 5.4 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and let $a, b \in A$. If $0 \sqsubseteq a$ and $0 \sqsubseteq b$ then $0 \sqsubseteq a \rightarrow b$.

Proof

$$[0 \to (a \to b)']' = [(a \to b) \to 0']'$$
by Lemma 2.5 (a)
$$= (0 \to a) \to (b \to 0')'$$
by (I)
$$= (0 \to a) \to (0 \to b')'$$
by Lemma 2.5 (a)
$$= (0 \to a) \to 0$$
since $0 \sqsubseteq b$
$$= 0' \to 0$$
by Lemma 5.3 since $0 \sqsubseteq a$
$$= 0.$$
by Lemma 2.4 (a)

Corollary 5.5 Let $\mathbf{A} \in \mathbf{I}_{2,0}$ and $a \in A$. If $a \supseteq 0$ then $a' \supseteq 0$.

Lemma 5.6 Let **A** be an $I_{2,0}$ -chain and let $a, b \in A$. Then $a' \to b' = b \to a$.

Proof Since **A** is a chain, we can assume that $b' \sqsubseteq a$ or $a \sqsubseteq b'$.

If $b' \sqsubseteq a$, $(b' \to a')' = b'$, then $b' \to a' = b$. Hence $b \to a = (b' \to a') \to a = [(a' \to b') \to (a' \to a)']'$, using (I). By Lemma 2.4 (d), $[(a' \to b') \to (a' \to a)']' = [(a' \to b') \to a']' = [[(a \to a') \to (b' \to a')$

If $a \sqsubseteq b'$ then we have $a' = (a \to b'')'' = a \to b$, and the rest of the argument is similar to the previous case.

Lemma 5.7 Let **A** be a $I_{2,0}$ -chain with $|A| \ge 2$ and let $a \in A$ such that $a \sqsubset 0$. Then

- (a) $0 \rightarrow a' = a'$
- (b) $0 \rightarrow a = a$
- (c) $(a \to a) \to a = a \to a$
- (d) $a \to a = a'$
- (e) $a \to a = a$
- (f) a = a'.

Proof

(a) Since $a \sqsubseteq 0$, we have that $a = (a \to 0')'$. Therefore, $a' = (a \to 0')'' = a \to 0' = 0 \to a'$ by Lemma 2.5 (b).

(b) Since $a \sqsubseteq 0$, we have

(5.2) $a = (a \to 0')'.$

Then we get

$$\begin{array}{rcl} (0 \rightarrow a) \rightarrow 0' &=& [(0 \rightarrow 0) \rightarrow (a \rightarrow 0')']' & \text{by (I)} \\ &=& [(0 \rightarrow 0) \rightarrow a]' & \text{by (5.2)} \\ &=& [0' \rightarrow a]' \\ &=& a' & \text{by lemma 2.4 (a)} \end{array}$$

Using Lemma 2.5 (b), we obtain

(5.3)
$$a' = 0 \to (0 \to a)'.$$

Since **A** is a chain, $0 \equiv 0 \rightarrow a$ or $0 \rightarrow a \equiv 0$. Suppose that $0 \equiv 0 \rightarrow a$. Then $(0 \rightarrow (0 \rightarrow a)')' = 0$ Therefore, by (5.3), $a = a'' = (0 \rightarrow (0 \rightarrow a)')' = 0$, a contradiction, since $a \neq 0$. Consequently, $0 \rightarrow a \equiv 0$. Hence, we have

$$0 \rightarrow a = ((0 \rightarrow a) \rightarrow 0')' \text{ since } 0 \rightarrow a \sqsubseteq 0$$

= $(0 \rightarrow (0 \rightarrow a)')'$ by lemma 2.5 (b)
= a'' by (5.3)
= a .

(c)

$$a \rightarrow a = (0 \rightarrow a) \rightarrow a \qquad \text{by item (b)}$$

$$= (a' \rightarrow 0') \rightarrow a \qquad \text{by lemma 2.5 (a)}$$

$$= [(a' \rightarrow a') \rightarrow (0' \rightarrow a)']' \qquad \text{by Lemma 5.6}$$

$$= [(a \rightarrow a) \rightarrow a']' \qquad \text{by lemma 2.5 (a)}$$

$$= [(a'' \rightarrow a) \rightarrow (a \rightarrow a')']' \qquad \text{by lemma 2.5 (a)}$$

$$= [(a'' \rightarrow a) \rightarrow (a \rightarrow a')']' \qquad \text{by (I)}$$

$$= (a \rightarrow a) \rightarrow (a'' \rightarrow a')'$$

$$= (a \rightarrow a) \rightarrow a'' \qquad \text{by lemma 2.5 (d)}$$

$$= (a \rightarrow a) \rightarrow a.$$

(d) Since **A** is a chain, $0 \to a' \sqsubseteq a$ or $a \sqsubseteq 0 \to a'$. First, we assume that $0 \to a' \sqsubseteq a$. Then

$$\begin{aligned} a \to a &= (a \to a) \to a & \text{by (c)} \\ &= (a' \to a') \to a & \text{by Lemma 5.6} \\ &= a' \to (a' \to a')' & \text{by Lemma 5.6} \\ &= (a \to 0) \to (a' \to a')' \\ &= [(a \to 0) \to (a' \to a')']'' \\ &= [(0 \to a') \to a']' & \text{using (I)} \\ &= 0 \to a' & \text{since } 0 \to a' \sqsubseteq a \\ &= a' & \text{using (a).} \end{aligned}$$

Next, we assume $a \sqsubseteq 0 \to a'$, i.e., $(a \to (0 \to a')')' = a$. Then, from (a), we have $a \to a = (a \to a'')'' = [(a \to (0 \to a')')']' = a'$.

- (e) Using the items (c), (d) and Lemma 2.5 (d), we have $a \to a = (a \to a) \to a = a' \to a = a$.
- (f) This follows immediately from the two preceding items.

Lemma 5.8 Let **A** be an $\mathbf{I}_{2,0}$ -chain with $|A| \ge 2$, and let $a, b \in A$. If $0 \sqsubseteq a$ and $b \sqsubset 0$ then $b \rightarrow a = b$ and $a \rightarrow b = b$.

Proof Since $0 \sqsubseteq a$ and $b \sqsubset 0$, we have that $(0 \to a')' = 0$ and $(b \to 0')' = b$. Therefore, using Lemma 5.6, $b = b'' = b' \to 0 = (b \to 0')'' \to (0 \to a')' = (b \to 0') \to (0 \to a')' = (0 \to b') \to (a \to 0')' = [(b' \to a) \to 0']'$. Hence,

(5.4) $b = [(b' \to a) \to 0']'.$

From the hypothesis and Lemma 5.7 (f), we have

$$(5.5) \ b' = b.$$

Suppose that $0 \sqsubseteq b' \to a$. Then $0 = [0 \to (b' \to a)']' = [(b' \to a) \to 0']'$ by Lemma 5.6, implying 0 = b, which is a contradiction in view of (5.4). Consequently, $b' \to a \sqsubseteq 0$, since **A** is a chain. Hence,

(5.6) $b' \to a = [(b' \to a) \to 0']'.$

From (5.4), (5.5) and (5.6) we conclude $b = b \rightarrow a$, proving the first half of the conclusion of the lemma. From

$$b = (b \to a')' \quad \text{since } b \sqsubseteq a, \text{ as } 0 \sqsubseteq a \text{ and } b \sqsubset 0$$
$$= (a'' \to b')' \quad \text{by Lemma 5.6}$$
$$= (a \to b')'$$
$$= (a \to b)' \quad \text{by (5.5)}$$

we conclude that $a \to b = b' = b$ in view of (5.5), completing the second half.

Definition 5.9 Let $\mathbf{A} = \langle A; \to, 0 \rangle$ be a finite $\mathbf{I}_{2,0}$ -chain. We let $A^+ := \{a \in A : a \sqsupset 0\}$ and $A^- := \{a \in A : a \sqsubset 0\}$. Observe that $A = A^+ \cup \{0\} \cup A^-$. Henceforth, without loss of generality, we will represent A = [-n, m] with $0 \le n, m \le |A| - 1$, such that

$$-n \sqsubset -n + 1 \sqsubset \ldots \sqsubset -1 \sqsubset 0 \sqsubset 1 \sqsubset 2 \sqsubset \ldots \sqsubset m.$$

Remark 5.10 In view of the above definition, we can use the functions * and p of Definition 4.2 as functions on the domain [-n, m] of **A** as well.

Now, we wish to prove that $\langle A; \rightarrow, 0 \rangle = \langle [-n; m]; \Rightarrow, 0 \rangle$. To achieve this, we need the following lemmas.

Lemma 5.11 Let $\mathbf{A} = \langle A; \rightarrow, 0 \rangle$ be a finite $\mathbf{I}_{2,0}$ -chain with $|A| \ge 2$. If $a \supseteq 0$ then a' = p(p(a)').

Proof By hypothesis we have that $a \supseteq 0$. Then $p(a) \supseteq 0$. Hence $0 \sqsubseteq p(a) \sqsubset a$. Then, by Lemma 5.2,

(5.7) $a' \sqsubseteq p(a)'$.

Since $a \equiv 0$, by Corollary 5.5, $a' \equiv 0$. Therefore, by (5.7),

(5.8) $0 \sqsubseteq p(a)'$.

If a' = p(a)' then a = p(a) and, consequently, a = -n, a contradiction, so $a' \sqsubset p(a)'$, and hence, $0 \sqsubseteq a' \sqsubseteq p(p(a)') \sqsubset p(a)'$. By lemma 5.2, $a \sqsupseteq [p(p(a)')]' \sqsupseteq p(a)$. Thus

(5.9) $[p(p(a)')]' \in \{a, p(a)\}.$

If [p(p(a)')]' = p(a), we have that p(p(a)') = [p(p(a)')]'' = p(a)', a contradiction, since $p(a)' \supseteq 0$ by (5.8). Therefore [p(p(a)')]' = a and therefore, p(p(a)') = a'.

Lemma 5.12 Let $\mathbf{A} = \langle A; \rightarrow, 0 \rangle$ be a finite $\mathbf{I}_{2,0}$ -chain. If $a \in A$ then $a^* = a'$.

Proof The statement $0' = m = 0^*$ follows from Lemma 5.1. If $a \sqsubset 0$ then a' = a by Lemma 5.7 (f), and $a = a^*$ by definition, implying $a = a^*$.

Now assume that $a \supseteq 0$. We will verify that $a' = a^*$ by induction on a. If a = 1, then, as $0' = 0^*$, we have, by Lemma 5.11, that $1' = p(p(1)') = p(0') = p(0^*) = p(p(1)^*) = 1^*$. The inductive hypothesis is that $p(a)' = p(a)^*$. Hence, we have, by Lemma 5.11, $a' = p(p(a)') = p(p(a)^*) = a^*$.

The following theorem shows that the general method described in Definition 4.2 essentially gives all finite $I_{2,0}$ -chains.

Theorem 5.13 Let **A** be a finite $\mathbf{I}_{2,0}$ -chain. Then $\mathbf{A} \cong \langle [-n,m]; \Rightarrow, 0 \rangle$ for some $0 \leq n, m \leq |A| - 1$.

Proof We will use the notation of Definition 5.9. Let $i, j \in A$. From Lemma 5.12, $i' = i^*$ and $j' = j^*$. It suffices to verify that

$$i \to j = \begin{cases} max(i',j) & \text{if } i,j \supseteq 0\\ min(i,j) & \text{otherwise} \end{cases}$$

with 0' = m. We consider the following cases:

• Case 1: j > 0.

We need the following subcases:

- Case 1.1: i > 0.

We make the following further subcases:

- * Case 1.1.1: $i' \ge j$. Since $i' \supseteq j$, we observe that
 - (5.10) $(j \to i'')' = j.$

Hence

$$i \rightarrow j = i \rightarrow (j \rightarrow i'')' \quad \text{by (5.10)}$$

$$= i \rightarrow (j \rightarrow i)'$$

$$= [(i \rightarrow j) \rightarrow i]' \quad \text{by Lemma 2.6 (4)}$$

$$= [(0 \rightarrow j) \rightarrow i]' \quad \text{by Lemma 2.6 (5)}$$

$$= [0' \rightarrow i]' \qquad \text{by Lemma 5.3 since } j \supseteq 0$$

$$= i' \qquad \text{by Lemma 2.4 (a)}$$

$$= max(i', j) \qquad \text{since } i' \supseteq j$$

* Case 1.1.2: i' < j. Since $i' \sqsubseteq j$, we have

(5.11)
$$(i' \to j')' = i'.$$

Therefore,

$$i \rightarrow j = i'' \rightarrow j$$

$$= (i' \rightarrow j')'' \rightarrow j \quad \text{by (5.11)}$$

$$= (i' \rightarrow j') \rightarrow j$$

$$= (i' \rightarrow 0') \rightarrow j \quad \text{by Lemma 2.6 (1)}$$

$$= (0 \rightarrow i) \rightarrow j \quad \text{by Lemma 2.5 (a)}$$

$$= 0' \rightarrow j \quad \text{by Lemma 5.3 since } i \sqsupseteq 0$$

$$= j \quad \text{by Lemma 2.4 (a)}$$

$$= max(i', j) \quad \text{since } i' \sqsubseteq j$$

- Case 1.2: i = 0.

Using Lemma 5.3 and Lemma 5.1, $0 \rightarrow j = 0' = max(0', j)$.

- Case 1.3: i < 0.

$$i \rightarrow j = (0 \rightarrow i) \rightarrow j \qquad \text{by Lemma 5.7 (b)}$$

$$= (i' \rightarrow 0') \rightarrow j \qquad \text{by Lemma 5.7 (f)}$$

$$= (i \rightarrow 0') \rightarrow j \qquad \text{by Lemma 5.7 (f)}$$

$$= [(j' \rightarrow i) \rightarrow (0' \rightarrow j)']' \qquad \text{by (I)}$$

$$= [(0 \rightarrow i) \rightarrow j']' \qquad \text{by Lemma 2.6 (5)}$$

$$= (i \rightarrow j')' \qquad \text{by Lemma 5.7 (b)}$$

$$= i \qquad \text{since } i \sqsubset j$$

$$= \min(i, j)$$

• Case 2: j < 0.

It is useful to consider the following subcases:

- Case 2.1: i > 0

$$i \rightarrow j = i \rightarrow j' \qquad \text{by Lemma 5.7 (f)}$$

$$= i \rightarrow (j \rightarrow i')'' \qquad \text{since } j \sqsubset i$$

$$= i \rightarrow (j \rightarrow i')$$

$$= j \rightarrow i' \qquad \text{by Lemma 2.6 (17)}$$

$$= (j \rightarrow i')''$$

$$= j' \qquad \text{since } j \sqsubset i$$

$$= j \qquad \text{by Lemma 5.7 (f)}$$

$$= \min(i, j)$$

- Case 2.2: i = 0.

$$i \rightarrow j = 0 \rightarrow j$$

= j by Lemma 5.7 (b)
= $min(i, j)$

- Case 2.3: i < 0.

* Case 2.3.1: $i \leq j$. As $i \sqsubseteq j$, we have

(5.12) $(i \to j')' = i.$

Observe

* Case 2.3.2: i > j. We have

i

(5.13) $(j \to i')' = j.$

as $j \sqsubseteq i$. Hence

$$i \rightarrow j = j' \rightarrow i' \quad \text{by Lemma 5.6}$$

= $j \rightarrow i' \quad \text{by Lemma 5.7 (f)}$
= $(j \rightarrow i')''$
= $j' \quad \text{by (5.13)}$
= $j \quad \text{by Lemma 5.7 (f)}$
= $min(j,i)$

• Case 3: j = 0.

- Case 3.1: $i \ge 0$. By Corollary 5.5, as $i \sqsupseteq 0$, we have that $i' = i \rightarrow 0 \sqsupseteq 0$. Hence $i \rightarrow 0 = i' = max(i', 0)$.

- Case 3.2: i < 0. We have that

$$i \rightarrow j = i \rightarrow 0$$

= i'
= i by Lemma 5.7 (f)
= min(i, j)

Hence $\mathbf{A} \cong \langle [-n;m]; \Rightarrow, 0 \rangle$.

The following theorem, our second main result, is now immediate from the preceding results.

Theorem 5.14 There are n non-isomorphic $I_{2,0}$ -chains of size n, for $n \in \mathbb{N}$.

A Appendix: Proofs

We would like to mention here that the identity: $x'' \approx x$ is used in these proofs frequently without explicit mention.

Proof of Lemma 2.6: Items (1) to (17) are proved in [3]. The proofs of (18) to (26) are given in [5]. Let $a, b, c, d \in A$.

(27)

$$(b \to c) \to a = [(b \to c) \to a]' \to [(b \to c) \to a]$$
by Lemma 2.4 (d)

$$= [(a' \to b) \to (c \to a)']'' \to [(b \to c) \to a]$$
from (I)

$$= [(a' \to b) \to (c \to a)'] \to [(b \to c) \to a]$$

(28)

$$\begin{split} \left[[0 \rightarrow (a \rightarrow b)'] \rightarrow (0 \rightarrow b')' \right]' &= \left[\{ 0 \rightarrow (a \rightarrow b)' \} \rightarrow (b \rightarrow 0')' \right]' & \text{by Lemma 2.5 (a)} \\ &= \left[(a \rightarrow b)' \rightarrow b \right] \rightarrow 0' & \text{by (I)} \\ &= 0 \rightarrow \left[(a \rightarrow b)' \rightarrow b \right]' & \text{by Lemma 2.5 (a)} \\ &= (a \rightarrow b) \rightarrow (0 \rightarrow b') & \text{by (10)} \\ &= 0 \rightarrow \left[(a \rightarrow b) \rightarrow b' \right] & \text{by (13)} \\ &= 0 \rightarrow \left[(a \rightarrow 0') \rightarrow b' \right] & \text{by (11)} \\ &= (a \rightarrow 0') \rightarrow (0 \rightarrow b') & \text{by (13)} \\ &= (0 \rightarrow a') \rightarrow (0 \rightarrow b') & \text{by Lemma 2.5 (a)} \\ &= a' \rightarrow (0 \rightarrow b') & \text{by (20)} \\ &= \left[(0 \rightarrow a) \rightarrow (b \rightarrow 0')' \right]' & \text{by (8)} \\ &= \left[(0 \rightarrow a) \rightarrow (b \rightarrow 0')' \right]' & \text{by Lemma 2.5 (a)} \\ &= (a \rightarrow b) \rightarrow 0' & \text{by (I)} \\ &= 0 \rightarrow (a \rightarrow b)' & \text{by Lemma 2.5 (a)} \\ \end{split}$$

(29)

$$0 \to [(a \to b) \to c] = 0 \to [(c' \to a) \to (b \to c)']' \text{ by (I)}$$
$$\sqsubseteq 0 \to (b \to c)'' \text{ by (28)}$$
$$= 0 \to (b \to c)$$

(30)

$$\begin{aligned} a' \to (b \to 0')' &= (a \to 0) \to (b \to 0')' \\ &= [\{(b \to 0') \to a\} \to \{0 \to (b \to 0')'\}']' & \text{by (I)} \\ &= [\{(b \to 0') \to a\} \to \{0 \to (0 \to b')'\}']' & \text{by Lemma 2.5 (a)} \\ &= [\{(b \to 0') \to a\} \to (0 \to b)']' & \text{by (9)} \\ &= [\{(0 \to b') \to a\} \to (0 \to b)']' & \text{by Lemma 2.5 (a)} \\ &= [[\{0 \to (0 \to b)'\} \to a] \to (0 \to b)']' & \text{by (9)} \\ &= [a \to (0 \to b)']' & \text{by (18)} \end{aligned}$$

(31)

$$[(0 \to a) \to b]' = [(b \to a) \to b]' \quad \text{by (5)}$$
$$= [b \to (a \to b)']'' \quad \text{by (4)}$$
$$= b \to (a \to b)'$$

(32)

$$[a \to (b \to 0')']' = [a \to (0 \to b')']'$$
by Lemma 2.5 (a)
$$= a' \to (b' \to 0')'$$
by (30)
$$= a' \to (0 \to b)'$$
by Lemma 2.5 (a)

(33)

$$b' \to a' = (b \to 0) \to a'$$

= $[(a \to b) \to (0 \to a')']'$ by (I)
= $[(a \to b) \to (a \to 0')']'$ by Lemma 2.5 (a)
= $(a \to b)' \to (0 \to a)'$ by (32) with $x = a \to b, y = a$

(34)

$$[(a \to b)' \to \{b \to (a \to b)'\}']' = [(a \to b)' \to b] \to (a \to b)' \quad \text{by (4)}$$
$$= (0 \to b) \to (a \to b)' \quad \text{by (5)}$$
$$= (a \to b)' \quad \text{by (3)}$$

(36)

$$\begin{array}{rcl} (0 \to a) \to b &=& (b \to a) \to b & \text{by (5)} \\ &=& [b \to (a \to b)']' & \text{by (4)} \\ &\sqsubseteq& (a \to b)' \to [b \to (a \to b)']' & \text{by (35) with } x = b, y = (a \to b)' \\ &=& [\{(a \to b)' \to b\} \to (a \to b)']' & \text{by (4)} \\ &=& [(0 \to b) \to (a \to b)']' & \text{by (5)} \\ &=& (a \to b)'' & \text{by (5)} \\ &=& (a \to b)'' & \text{by (3)} \\ &=& a \to b & \text{since } x'' \approx x \end{array}$$

$$\begin{split} [\{a \to (b \to a)'\} \to a'']' &= [\{a \to (b \to a)'\} \to a]' \\ &= [\{0 \to (b \to a)'\} \to a]' \quad \text{by (5)} \\ &= [\{(b \to a) \to 0'\} \to a]' \quad \text{by Lemma 2.5 (a)} \\ &= [\{(b \to a) \to a'\} \to a]' \quad \text{by (1)} \\ &= [\{(b \to 0') \to a'\} \to a]' \quad \text{by (1)} \\ &= [\{(b \to 0') \to 0'\} \to a]' \quad \text{by (1)} \\ &= [\{(b \to 0'') \to 0'\} \to a]' \quad \text{by (1)} \\ &= [\{(b \to 0'') \to 0'\} \to a]' \quad \text{by (1)} \\ &= [\{(b \to 0) \to 0'\} \to a]' \\ &= [(b' \to 0') \to a]' \\ &= [(b' \to 0') \to a]' \quad \text{by Lemma 2.5 (a)} \\ &= [(a \to b) \to a]' \quad \text{by (5)} \\ &= a \to (b \to a)' \quad \text{by (4)} \end{split}$$

Proof of Lemma 3.3

(1) Observe that by Lemma 2.5 (a), Lemma 2.6 (1) and the hypothesis we have that $(0 \rightarrow a') \rightarrow b = (a \rightarrow 0') \rightarrow b = (a \rightarrow b') \rightarrow b = (a \rightarrow b')'' \rightarrow b = a' \rightarrow b$.

(37)

$$b \to a' = [(0 \to a') \to b] \to a' \text{ by Lemma 2.6 (18)}$$
$$= (a' \to b) \to a' \text{ from (1)}$$
$$= (0 \to b) \to a' \text{ by Lemma 2.6 (5).}$$

(3)

$$b \rightarrow a' = (0 \rightarrow b) \rightarrow a' \qquad \text{from (2)}$$

= $(0 \rightarrow b) \rightarrow (a \rightarrow b')''$ by hypothesis
= $(0 \rightarrow b) \rightarrow (a \rightarrow b')$
= $(0 \rightarrow b'') \rightarrow (a \rightarrow b')$
= $a \rightarrow b'$ by Lemma 2.6 (2)
= $(a \rightarrow b')''$
= a' by hypothesis

(4)

$$\begin{array}{rcl} 0 \rightarrow (a' \rightarrow b) &=& a' \rightarrow (0 \rightarrow b) & \text{by Lemma 2.6 (13)} \\ &=& 0 \rightarrow (a \rightarrow b')' & \text{by Lemma 2.6 (10)} \\ &=& 0 \rightarrow a & \text{by hypothesis} \end{array}$$

(5) By hypothesis and (I) we have that $(d \to a) \to b' = [(b \to d) \to (a \to b')']' = [(b \to d) \to a]'$.

(6)

$$[\{d \to (0 \to b')\} \to a]' = (a' \to d) \to [(0 \to b') \to a]' \text{ by (I)}$$

$$= (a' \to d) \to [(a \to b') \to a]' \text{ by Lemma 2.6 (5)}$$

$$= (a' \to d) \to (a' \to a)' \text{ by hypothesis}$$

$$= (a' \to d) \to (a' \to a)' \text{ by Lemma 2.4 (d)}$$

$$= (a' \to d) \to (0' \to a)' \text{ by Lemma 2.4 (a)}$$

$$= [(d \to 0') \to a]' \text{ by (I)}$$

$$= (0 \to d) \to a' \text{ by Lemma 2.6 (11)}$$

(7)

$$\begin{aligned} a \to [(a' \to d) \to \{(0 \to a) \to b'\}] &= a \to [(a' \to d) \to \{(b \to 0) \to (a \to b')'\}'] \\ & \text{by (I)} \\ &= a \to [(a' \to d) \to \{(b \to 0) \to a\}'] \\ & \text{by hypothesis} \\ &= a \to [\{d \to (b \to 0)\} \to a]' \\ & \text{by (I)} \\ &= [[d \to (0 \to (b \to 0))] \to a]' \\ & \text{by Lemma 2.6 (22) with } x = d, y = b \to 0, z = a \\ &= [[d \to (0 \to b')] \to a]' \\ &= (0 \to d) \to a' \\ & \text{by (6)} \end{aligned}$$

(8)

(9)

$$\begin{bmatrix} 0 \to (b \to d) \end{bmatrix} \to a = \begin{bmatrix} (a' \to 0) \to ((b \to d) \to a)' \end{bmatrix}' \text{ by (I)}$$

$$= \begin{bmatrix} a \to ((b \to d) \to a)' \end{bmatrix}'$$

$$= \begin{bmatrix} a \to ((d \to a) \to b') \end{bmatrix}' \text{ by (5)}$$

$$= \begin{bmatrix} a \to (d \to a)' \end{bmatrix}'$$

$$= \begin{bmatrix} a \to (d \to a)' \end{bmatrix}'$$

$$= \begin{bmatrix} a \to (d \to a)' \end{bmatrix}'$$

$$= \begin{bmatrix} a \to (d \to a) \\ 0 \to a \end{bmatrix}$$

$$= \begin{bmatrix} a \to d \\ 0 \to d \end{bmatrix} \to a$$

$$= \begin{bmatrix} 0 \to d \\ 0 \to a \end{bmatrix}$$

$$= \begin{bmatrix} 0 \to d \\ 0 \to a \end{bmatrix}$$

$$= \begin{bmatrix} 0 \to d \\ 0 \to a \end{bmatrix}$$

$$= \begin{bmatrix} 0 \to d \\ 0 \to a \end{bmatrix}$$

(10)

$$(b \to (a \to d)) \to a = [(a' \to b) \to \{(a \to d) \to a\}']' \quad \text{by (I)}$$

$$= [(a' \to b) \to \{(0 \to d) \to a\}']' \quad \text{by Lemma 2.6 (5)}$$

$$= [b \to (0 \to d)] \to a \qquad \qquad \text{by (I)}$$

$$= [0 \to (b \to d)] \to a \qquad \qquad \text{by Lemma 2.6 (13)}$$

$$= (0 \to d) \to a \qquad \qquad \qquad \text{by (9)}$$

(11)

$$b \to (0 \to a') = (0 \to b) \to (0 \to a') \qquad \text{by Lemma 2.6 (20)}$$

$$= 0 \to [(0 \to b) \to a'] \qquad \text{by Lemma 2.6 (13)}$$

$$= 0 \to [(a \to 0) \to (b \to a')']' \qquad \text{by (I)}$$

$$= 0 \to [a' \to (b \to a')']' \qquad \qquad \text{by (3)}$$

$$= 0 \to (a' \to a)' \qquad \qquad \qquad \text{by Lemma 2.4 (d)}$$

(12) From (I) and by hypothesis we have that $[(d \to a) \to b']' = (b \to d) \to (a \to b')' = (b \to d) \to a$.

$$\begin{array}{rcl} a' \rightarrow b &=& (a \rightarrow b') \rightarrow b & \text{by hypothesis} \\ &=& [(b' \rightarrow a) \rightarrow (b' \rightarrow b)']' & \text{using (I)} \\ &=& [(b' \rightarrow a) \rightarrow b']' & \text{by Lemma 2.4 (d)} \\ &=& b' \rightarrow (a \rightarrow b')' & \text{by Lemma 2.6 (4)} \\ &=& b' \rightarrow a & \text{by hypothesis} \end{array}$$

(14)

(15)

$$[(0 \to a') \to b]' = [(a \to 0') \to b]' \text{ by Lemma 2.5 (a)}$$
$$= (0 \to a) \to b' \text{ by Lemma 2.6 (11)}$$

(16)

$$\begin{aligned} (a' \to b)' &= [(0 \to a') \to b]' & \text{by (1)} \\ &= (0 \to a) \to b' & \text{by (15)} \end{aligned}$$

(17)

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