

Lost in translation: unknowable propositions in probabilistic frameworks

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Abstract Some propositions are *structurally unknowable* for certain agents. Let me call them ‘Moorean propositions’. The structural unknowability of Moorean propositions is normally taken to pave the way towards proving a familiar paradox from epistemic logic—the so-called ‘Knowability Paradox’, or ‘Fitch’s Paradox’—which purports to show that if all truths are knowable, then all truths are in fact known. The present paper explores how to translate Moorean statements into a probabilistic language. A successful translation should enable us to derive a version of Fitch’s Paradox in a probabilistic setting. I offer a suitable schematic form for probabilistic Moorean propositions, as well as a concomitant proof of a probabilistic Knowability Paradox. Moreover, I argue that traditional candidates to play the role of probabilistic Moorean propositions will not do. In particular, we can show that violations of the so-called ‘Reflection Principle’ in probability (as discussed for instance by Bas van Fraassen) need not yield structurally unknowable propositions. Among other things, this should lead us to question whether violating the Reflection Principle actually amounts to a clear case of epistemic irrationality, as it is often assumed. This result challenges the importance of the principle as a tool to assess both synchronic and diachronic rationality—a topic which is largely independent of Fitch’s Paradox—from a somewhat unexpected source.

Keywords Knowability · Fitch’s Paradox · Moore · Reflection Principle

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20 1 Introduction

21 There is a familiar paradox from epistemic logic—the so-called ‘Knowability Para-
 22 dox’, or ‘Fitch’s Paradox’—which purports to show that if all truths are knowable, then
 23 all truths are in fact known; to put it differently, knowability (or weak verificationism)
 24 collapses with omniscience (or strong verificationism). This result is paradoxical if we
 25 think, as it seems natural to do, that there are indeed truths that nobody knows, and if
 26 we also think that the antecedent, the Knowability Principle (‘all truths are knowable’)
 27 is a substantive thesis whose truth or falsity should not be decided *a priori*.¹

28 Fitch’s paradox is intimately tied to so-called ‘Moore’s paradox’.² Moore calls our
 29 attention to the fact that no one can assert a statement such as ‘here is a rabbit, but I
 30 don’t believe it’, without causing perplexity in the audience—although, *sensu stricto*,
 31 we do not have a formal contradiction, at least not yet. This situation is even clearer if
 32 we replace belief by knowledge. A standard way of putting Fitch’s argument in motion
 33 is precisely by noticing that Moorean sentences of the type ‘*p*, but I don’t know that
 34 *p*’, if true, can’t themselves be known to be true: they are structurally un-knowable.³

35 Here is a brief reconstruction of Fitch’s alleged ‘modal collapse’ between actual
 36 and possible knowledge:

- | | | |
|----|---|---|
| 37 | 1. $K(\varphi \wedge \neg K\varphi)$ | [assumption] |
| 38 | 2. $K\varphi \wedge K\neg K\varphi$ | [distribution of K over conjunction] |
| 39 | 3. $\neg K\varphi$ | [from 2, by the factivity of K] |
| 40 | 4. \perp | [from 2 and 3, by propositional logic] |
| 41 | 5. $\vdash \neg K(\varphi \wedge \neg K\varphi)$ | [from 1–4] |
| 42 | 6. $\vdash \Box\neg K(\varphi \wedge \neg K\varphi)$ | [from 5, because theorems are
43 necessary] |
| 44 | 7. $\vdash \forall\psi(\psi \rightarrow \Diamond K\psi)$ | [Knowability Principle: every truth is
45 knowable] ⁴ |
| 46 | 8. $\vdash (\varphi \wedge \neg K\varphi) \rightarrow \Diamond K(\varphi \wedge \neg K\varphi)$ | [instance of 7] |
| 47 | 9. $\vdash \neg\Diamond K(\varphi \wedge \neg K\varphi)$ | [from 6, by definition of alethic
48 modalities] |
| 49 | 10. $\vdash \neg(\varphi \wedge \neg K\varphi)$ | [from 8 and 9] |

¹ The paradox appeared in press for the first time in Fitch (1963). In Salerno (2009) we can find a detailed account of the story of the paradox.

² Cf. Moore (1993).

³ A word of caution. It is not completely obvious that the epistemic version of Moore’s claim preserves the most interesting traits of Moore’s original paradox without turning it into something different. I will comment very briefly on this point below.

⁴ Arguably, a plausible version of this principle may require that the ‘ K ’ operator be read as ‘someone, at some time, knows that’. However, we may well adapt the Knowability Principle so that it could be represented within models that deal with the knowledge of a single (ideally rational) agent. On the other hand, some attempts to solve the paradox would contend that (7) does not express the intuitive idea of the knowability of truth—say, because it fails to include an actuality operator [as in Edgington’s proposal; cf. Edgington (1985)] or because, as it stands, (7) hides the relevant quantifiers, which should be understood as modal indexicals [as in Kvanvig’s account; cf. Kvanvig (2006)]. For the most part, in this paper I will not be concerned with possible ways to block the paradox, although I will have something to say about this in the last section.

- 50 11. $\vdash \forall \psi (\psi \rightarrow K \psi)$ [generalization from 10]
 51 12. $\vdash \forall \psi (\psi \rightarrow \Diamond K \psi) \rightarrow \forall \psi (\psi \rightarrow K \psi)$ [from 7–11]⁵

52 Given that truth implies possibility, the converse holds as well, so we can strengthen
 53 the result to a biconditional:

- 54 13. $\vdash \forall \psi (\psi \rightarrow \Diamond K \psi) \leftrightarrow \forall \psi (\psi \rightarrow K \psi)$ [from 12]

55 Moreover, if the Knowability Principle is accepted, we obtain an even stronger con-
 56 clusion:

- 57 14. $\vdash \forall \psi (\psi \rightarrow \Diamond K \psi) \wedge \forall \psi (\psi \rightarrow K \psi)$ [from 7 and 11]
 58 15. $\vdash \forall \psi (\psi \rightarrow (\Diamond K \psi \leftrightarrow K \psi))$ [from 14]

59 For the most part, what I have to say applies indistinctly to sentences (or statement)
 60 and to propositions; if ‘ φ ’ is a sentence, the proposition expressed by ‘ φ ’ is the set
 61 of all and only those possible worlds in which ‘ φ ’ comes out true (in a given inter-
 62 pretation). As is customary, I will represent propositions by means of square brackets
 63 on the corresponding sentences. Moreover, I will speak loosely of agents knowing or
 64 believing sentences, even though this is not the usual way to go. The reader can just
 65 take it as an abbreviation for ‘knowing or believing that the sentence is true’.

66 Let me introduce a bit of terminology. By a ‘Moorean statement’, or ‘Moorean sen-
 67 tence’, I mean any statement of the form ‘ $\varphi \wedge \neg \Delta \varphi$ ’, where ‘ φ ’ replaces any sentence,
 68 and ‘ Δ ’ is some epistemic or doxastic operator in a wide sense; *mutatis mutandis* for
 69 ‘Moorean proposition’. Notice that line (6) from the previous argument states that
 70 epistemic Moorean sentences cannot be known.

71 It should be noted that the connection between Moore’s and Fitch’s paradoxes works
 72 only to the extent that the belief operator can be credited with the satisfaction of the
 73 right combination of principles, such as distribution over conjunction ($\vdash B(\varphi \wedge \psi) \rightarrow$
 74 $(B\varphi \wedge B\psi)$) and moderate factivity ($\vdash B\neg B\varphi \rightarrow \neg B\varphi$); alternatively, we can avoid
 75 moderate factivity but demand that distribution over conjunction be strengthened to
 76 a biconditional, while also demanding doxastic transparency ($\vdash B\varphi \rightarrow BB\varphi$) and
 77 doxastic consistency ($\vdash \neg B(\varphi \wedge \neg \varphi)$); or perhaps we would rather have distribution
 78 over conjunction plus extended doxastic consistency (‘if $B\varphi$ and $B\psi$, then $(B\varphi \wedge \psi) \vdash$
 79 \perp ’). We might also need to claim that ‘sincere assertion bestows belief’ (as in [Tennant](#)
 80 [\(1997\)](#)), insofar as it is the mere *assertion* of a (doxastic) Moorean statement which
 81 suffices to cause trouble. I am ready to concede that at least some such principles
 82 are fine as far as *rational* belief is concerned, and hence that we can indeed draw
 83 an analogy between Moore’s paradox and step (6) from the Knowability argument.⁶
 84 However, it might be argued that the original Moorean statement need not be about
 85 rational belief at all, and hence that no interesting connection can be made with Fitch’s
 86 Paradox.⁷ I will not try to settle this debate here. If the reader feels inclined to reject

⁵ To avoid quantification over propositions, one might treat (7) and (12) as sentence schemata and thus omit propositional quantifiers; the same applies to (13)–(15) below (thanks to an anonymous referee for pressing this point).

⁶ For sympathetic approaches to the idea that there is an interesting connection between the two paradoxes cf. [Tennant \(1997\)](#), [Linsky \(2009\)](#), or [van Benthem \(2004\)](#), among other authors.

⁷ Cf. [Kvanvig \(2006\)](#).

87 the connection, then just assume that I am talking about Moorean* sentences and
 88 Moore's Paradox*, where Moore's Paradox* enables us to show the inconsistency of
 89 ' $B(\varphi \wedge \neg B\varphi)$ '.

90 I will also say that a 'quasi-Moorean statement' captures the essence of Moorean
 91 sentences in a probabilistic framework. Here are some possible examples, where ' $[\varphi]$ '
 92 is the proposition expressed by ' φ ', ' P ' is a subjective, or perhaps an evidential,
 93 probability function, and r and s are real numbers in $[0,1]$:

$$\begin{aligned}
 94 \quad & P([\varphi] \mid P([\varphi]) = r) \neq r \\
 95 \quad & P([\varphi]) = r \wedge P(P([\varphi]) = r) \neq 1 \\
 96 \quad & \varphi \wedge P([\varphi]) \neq 1 \\
 97 \quad & \varphi \wedge P([\varphi]) = 0 \\
 98 \quad & \varphi \wedge P([\varphi]) < r \text{ for some acceptance threshold } r
 \end{aligned}$$

99 I will distinguish between *potential* and *genuine* quasi-Moorean statements; the
 100 second ones are the successful candidates. We know we have found a good transla-
 101 tion of a Moorean sentence into a probabilistic framework if we are able to preserve
 102 essential features of Moorean statements in a probabilistic realm. I will call it 'the
 103 symmetry criterion'. Among other things, our candidate should preserve the ability to
 104 trigger a Fitch-like paradox.

105 In this paper I will explore the structure of several quasi-Moorean statements and use
 106 them to reconstruct a concomitant version of Fitch's paradox so as to get 'probabilistic
 107 un-knowability', so to speak. To carry out this project I will rely on a Kripke model
 108 enriched with probabilities. Reflecting on which candidates work, and which ones do
 109 not work, will prove useful to draw a number of morals that go well beyond the realm
 110 of epistemic or doxastic paradoxes. Ultimately, my purpose is to use this discussion
 111 to shed some light on the adequacy of certain epistemic principles. In particular, I will
 112 examine the putative existence of an appropriate link between the so-called Reflection
 113 Principle in probability and Moorean statements; this, in turn, will raise some doubts
 114 on the thought that satisfying the Reflection Principle is mandated by rationality. In
 115 addition, the project will highlight the convenience of adopting a possible refinement
 116 of the formalism, within which both strands of the paradox get a unified solution, in
 117 agreement with the symmetry criterion. I will not elaborate on the details here, as I
 118 have already presented the refinement somewhere else.⁸ I believe the proposed model
 119 is adequate on independent grounds; the fact that it allows for a unified solution should
 120 reinforce our confidence in its adequacy.

121 What is the significance of Fitch's paradox? As is well known, there is no consensus
 122 on how to answer this question. To begin with, we can wonder in what sense we have
 123 a *bona fide* paradox, and not just a *reductio* of the Knowability Principle. Given that at
 124 least some versions of semantic anti-realism might want to embrace the claim that all
 125 truths are in principle knowable, Fitch's argument has sometimes been interpreted as a
 126 refutation of certain types of anti-realism (Hart and McGinn (1976)). Alternatively, it
 127 has been contended that Fitch's strategy is not problematic for the anti-realist once the
 128 Knowability Principle is correctly formulated (Edgington (1985), Edgington (2010))

⁸ Cresto (2012).

129 or suitably restricted (Tennant (1997)), or once we realize that the anti-realist should
 130 be committed to the use of intuitionistic logic (Williamson (1982)). Yet other authors
 131 have suggested that Fitch’s result is hard to swallow even for those who have no
 132 interest in semantic anti-realism whatsoever (Kvanvig (2006)). In this paper I will
 133 remain neutral on this controversy. Recall, moreover, that even though I believe a
 134 probabilistic version of the paradox can have an interest in itself, my main goal will
 135 be to use Fitch’s paradox as a litmus test that will help us assess candidates for quasi-
 136 Moorean statements, following the symmetry criterion. This task is compatible with all
 137 major interpretations I have just mentioned. Whatever it is that is deemed problematic
 138 about the Knowability paradox, the problem should be inherited by a probabilistic
 139 setting; alternatively, if all Fitch’s argument achieves is a *reductio* of verificationism,
 140 then we should find an analogous *reductio* within a probabilistic realm. (Presumably,
 141 within a probabilistic setting the verificationist should commit herself to a probabilistic
 142 knowability principle: true propositions should be able to have maximum evidential
 143 probability.)

144 2 A probabilistic setting

145 Consider a Kripke structure $S = \langle W, R, P_{prior}, v \rangle$ for a single agent, where ‘ W ’ is
 146 a countable set of possible worlds, ‘ R ’ is a suitable epistemic accessibility relation
 147 among worlds, ‘ P_{prior} ’ is a finitely additive prior probability function on subsets of
 148 W , and ‘ v ’ is a valuation function for the sentences of a suitably regimented lan-
 149 guage L . We shall assume regularity, *i.e.*, for any $A \subseteq W$, $P_{prior}(A) = 0$ iff $A = \emptyset$.
 150 This structure allows us to represent both knowledge and higher-order probability
 151 attributions. A similar account has been proposed by Timothy Williamson in recent
 152 years,⁹ and is also reminiscent of other well known proposals that combine probabili-
 153 ties with epistemic operators, such as Halpern (2003).¹⁰ Within Williamson’s original
 154 framework, ‘ P_{prior} ’ is meant to capture the intrinsic plausibility of worlds before the
 155 evidence comes in, but we need not commit ourselves to this particular interpretation;
 156 we can simply conceive of it as embodying the priors of the agent, or perhaps the
 157 priors ascribed to the agent by the theoretician—the one who attempts to make both
 158 knowledge and probability attributions to the agent, from a third person point of view.

159 Within this setting, define ‘ $R(w)$ ’, for any world w , as the strongest proposition
 160 known by the agent in w :

$$161 \quad R(w) = \{x : wRx\}$$

162 We will assume that R is reflexive. This guarantees that $R(w)$ is not empty, for any
 163 w ; it also guarantees the factivity of knowledge, as is well known.

164 Define next the evidential probability of any proposition ‘ $[\varphi]$ ’ in a given world w ,
 165 for any w , as the prior probability of ‘ $[\varphi]$ ’ conditional on the strongest proposition
 166 known by the agent in that world, *i.e.*:

⁹ Williamson (2014).

¹⁰ Halpern (2003), chapter 7.

167

$$P_w([\varphi]) = P_{prior}([\varphi] \mid R(w))$$

168 where ‘ φ ’ is any sentence of L . Recall that $R(w)$ is never empty and that we demanded
169 regularity; hence (unconditional) evidential probabilities are always well-defined.

170 We can also profit from Williamson’s device to refer to higher order probabilities.
171 We shall say that proposition ‘ $[P([\varphi]) = r]$ ’, which tells us that the probability of
172 ‘ $[\varphi]$ ’ is r , is the set of all worlds in which the evidential probability of ‘ $[\varphi]$ ’ is r (for
173 some r in $[0, 1]$):

174

$$[P([\varphi]) = r] = \{w : P_w([\varphi]) = r\}$$

175 Then propositions such as ‘ $[P([\varphi]) = r]$ ’ can be plugged in as further arguments of
176 the prior probability function of the model, thereby obtaining what could be understood
177 (arguably) as a second order probability.¹¹

178 Notice that expressions such as ‘ $P_w([\varphi]) = r$ ’ or ‘ $P_{prior}([\varphi]) = s$ ’ are metalinguistic,
179 and hence they do not belong to the object language. For the most part, here I
180 will leave the mechanism to build probabilistic sentences of L undetermined, *i.e.*, I will
181 not be explicit as to how to build sentences of L that encode the probabilistic commitments
182 of the agent. When needed, I will just underline the relevant proposition and use
183 the underlined expression as a shortcut for some sentence of L that expresses exactly
184 that proposition. Notice that, if it is true that $[\varphi]$ ’s evidential probability in world w is
185 r , then w belongs to the set of worlds picked out by proposition ‘ $[P([\varphi]) = r]$ ’, and
186 hence any sentence that expresses exactly that proposition will be true in w . In other
187 words, we have:

188

$$P_w([\varphi]) = r \text{ iff } S, w \models \underline{[P([\varphi]) = r]}$$

189 To illustrate briefly how the model works, consider the following toy example.
190 Suppose $W = \{w, x, y\}$, and suppose R is as shown in Fig. 1.

191 Here and elsewhere I use capital letters for sets of worlds, or propositions; when
192 needed, I will also keep on using sentences of L between square brackets. If we
193 assume a uniform prior probability function, we obtain:

194

$$P_{prior}(A) = 2/3$$

195

$$P_w(A) = 1/2$$

196

$$[P(A) = 1/2] = \{w\}$$

197

$$P_w([P(A) = 1/2]) = 1/2$$

198 In other words, A ’s prior probability is $2/3$, while its evidential probability in w is
199 just $1/2$. Moreover, w is the only world in which the evidential probability for A is
200 $1/2$. Hence the (second order) evidential probability in w of the proposition stating
201 that A ’s probability is $1/2$ is, again, $1/2$.

¹¹ In Sect. 9 I will address some worries on whether this analysis captures what we intuitively demand from a second order probability.

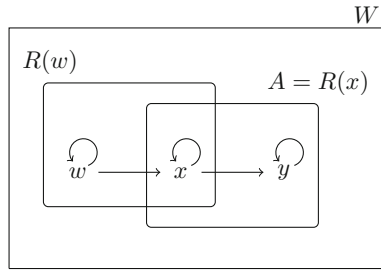


Fig. 1 A toy example for evidential probabilities

3 Quasi-Moorean statements

Are there any obvious candidates to formulate quasi-Moorean statements within framework S ? What we might call the ‘natural’ proposal asserts the conjunction of ‘ φ ’ and a (probabilistic) statement telling us that the probability of ‘ $[\varphi]$ ’ is less than 1:

$$\varphi \wedge [P([\varphi]) < 1]$$

We will examine this option with care in further sections. But first, we have some work to do. Interestingly, one of the most detailed discussion of the ‘natural’ candidate one can find in the literature seeks to establish a strong connection between such a statement and the so-called *Reflection Principle* in probability.¹² Moreover, the connection is meant to secure a possible line of defense for the Reflection Principle: violating the Reflection Principle would be analogous to committing Moore’s paradox; this is in particular van Fraassen’s position in van Fraassen (1995).

How is the connection between the Reflection Principle and Moore’s paradox supposed to go, exactly? One possibility is to contend that a sentence that *negates* the Reflection Principle is itself a genuine quasi-Moorean statement. This is not exactly van Fraassen’s position, but a close relative. A second, more cautious attitude would be to contend that the negation of the Reflection Principle mimics Moore’s paradox because this negation is entailed by the assumption that a statement such as ‘ $\varphi \wedge [P([\varphi]) < 1]$ ’ has itself maximum probability; this is actually the perspective endorsed by van Fraassen.¹³

In what follows I will begin by showing that the negation of the Reflection Principle is not a genuine quasi-Moorean statement. This result will clear the ground so that we can re-direct our efforts to more promising alternatives. In any case, a critical examination of why this identification fails will prove to be rewarding. I will discuss

¹² The Reflection Principle discussed in this section is not to be confused with ‘epistemic reflexivity’, most often referred to as ‘epistemic transparency’ or ‘the KK Principle’ (for any proposition φ : $\vdash K\varphi \rightarrow KK\varphi$). It is interesting to explore how the two senses of reflection interact with each other; I will take up this topic on board explicitly in further sections.

¹³ To wit: Assume $P([\varphi] \cap P([\varphi]) < 1) = 1$. Then $P([\varphi] \mid P([\varphi]) < 1)P(P([\varphi]) < 1) = 1$, which means that both factors are 1. As we will see, ‘ $P([\varphi] \mid P([\varphi]) < 1) = 1$ ’ is a special case of [RP Failure], as will be stated below.

227 more relaxed connections between the Reflection Principle and Moore's paradox in
228 later sections.

229 Some variants of the Reflection Principle are meant to refer to personal probabilities
230 ties; others seek to connect personal probabilities with chances, as in David Lewis's
231 *Principal Principle*. A possible version for evidential probabilities, within structure
232 S , may go as follows:

233 $P_w([\varphi] \mid [P([\varphi]) = r]) = r$, for all w in which the conditional probability **[RP]**
234 is defined
235

236 (*i.e.*, for all w such that $P_w([P([\varphi]) = r]) > 0$). That is to say, the evidential prob-
237 ability of $[\varphi]$ (in a particular world w), given that the probability of $[\varphi]$ is r , is also
238 r . One of the main reasons why [RP] has attracted so much interest in the litera-
239 ture is its potential connection with discussions on diachronic rationality; we might
240 acknowledge such a connection when we focus on versions of the principle in which
241 we conditionalize on propositions that announce the probabilities held by the agent
242 at later times, as in ' $P_{t_0}(A \mid [P_{t_1}(A) = r]) = r$ '.¹⁴ In addition, we might endorse a
243 version that makes room for vague probabilities, or vague partial beliefs. In this spirit,
244 van Fraassen's *General Reflection Principle* actually goes like this:

245 My current opinion about event E must lie in the range spanned by the possible
246 opinions I may come to have about E at later time t , as far as my present opinion
247 is concerned. (van Fraassen (1995), p. 16)¹⁵

248 In the rest of the paper I will focus exclusively on the synchronic case, and I will keep
249 on working with precise real numbers, for the sake of simplicity. In further sections,
250 however, I will seek to establish connections with other senses of vagueness.

251 Consider, then:

252 $P_w([\varphi] \mid [P([\varphi]) = r]) \neq r$ for some world w . **[RP Failure]**

253 I hope to show that, in spite of its initial plausibility, [RP Failure] is not a genuine
254 example of a quasi-Moorean statement.

255 Notice that, within setting S , the validity of [RP] depends on the structure of R . We
256 have a straightforward counterexample in Fig. 2:

257 For simplicity, let us assume once again a uniform prior probability distribution; we
258 then obtain:

259 $P_w(A) = 0$
260 $P_x(A) = 1/2$
261 $P_y(A) = 1/2$

¹⁴ Actually, van Fraassen (1995) defends [RP] as a modest constraint on diachronic rationality, as opposed to full-fledged Bayesian conditionalization.

¹⁵ Pace van Fraassen, it can be argued that the range of possible opinions I may come to have about E at a later time does not stand for a vague probability, but for a range of possible sharp probability assignments (thanks to an anonymous referee for pressing this point).

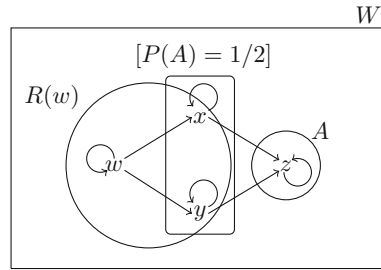


Fig. 2 A model for [RP Failure]

262 $P_w([P(A) = 1/2]) = 2/3$
 263 $P_w(A \mid [P(A) = 1/2]) = 0$

264 As we can see, [RP] has been just violated.

265 As a matter of fact, the semantics of our system guarantee that [RP] holds iff R
 266 is an equivalence relation. For an informal account of why this is so, notice that, by
 267 definition, the evidential probability of any proposition A , in some world w , refers to
 268 the probability of A given what the agent knows to be the case in that very world w .
 269 Therefore, we should expect that [RP] be satisfied when what the agent knows in w
 270 (i.e., ‘ $R(w)$ ’) does not have any chance of affecting her confidence in the proposition
 271 on which she is conditionalizing. This means, in turn, that the strongest proposition
 272 the agent knows in w is included in the proposition stating that A ’s probability is (say)
 273 r . Now, if R is an equivalence relation, a (second-order) probability on ‘ $[P(A) = r]$ ’
 274 will always be 0 or 1. So, if R is an equivalence relation, [RP] is either undefined
 275 or satisfied. As it turns out, this result can be strengthened to a biconditional. More
 276 precisely:

277 **Proposition 1** For any reflexive structure $S = \langle W, R, P_{prior}, v \rangle$: R is an equivalence
 278 relation iff for any world w , any proposition A and any r in $[0,1]$, $P_w(A \mid [P(A) =$
 279 $r]) = r$, if it is defined.¹⁶

280 *Proof* Left to right: By definition, $P_w(A \mid [P(A) = r]) = P_{prior}(A \cap \{w : P_w(A) =$
 281 $r\} \cap R(w)) / P_{prior}(\{w : P_w(A) = r\} \cap R(w))$. Assume R is an equivalence relation.
 282 Then R partitions the domain in such a way that, for all $y \in R(w)$, $R(w) = R(y)$.
 283 Hence for any A , if $y \in R(w)$, $P_y(A) = P_w(A)$. This amounts to saying that, for
 284 any particular $r \in [0, 1]$: either $P_y(A) = r$ for all $y \in R(w)$, or $P_y(A) \neq r$, for all
 285 $y \in R(w)$. Hence either $R(w) \subseteq \{w : P_w(A) = r\}$, or $R(w) \cap \{w : P_w(A) = r\} = \emptyset$.
 286 In the last case [RP] is undefined. If [RP] is not undefined, by contrast, $P_w[P(A) =$
 287 $r] = 1$. Our formula then reduces to $P_{prior}(A \cap R(w)) / P_{prior}(R(w)) = P_w(A) = r$.

288 Right to left: R is reflexive, by assumption. We will prove that it is also symmetric
 289 and transitive.¹⁷ To show transitivity, assume there are w, x and y in W such that wRx

¹⁶ This proposition has already been proven by Williamson (2014). By a ‘reflexive structure’ I mean any S with a reflexive R ; recall that we need to assume the reflexivity of R anyway in order to account for the factivity of knowledge.

¹⁷ For this part of the proof I follow closely Proposition 1 in Williamson (2014).

290 and xRy . As xRy , we have $P_x(\{y\}) = b > 0$. Hence $x \in [P(\{y\}) = b]$. Moreover, as
 291 $x \in R(w)$, we can guarantee that $P_w(\{y\} \mid [P(\{y\}) = b]) = P_{prior}(\{y\} \cap [P(\{y\}) =$
 292 $b]) \cap R(w) / P_{prior}([P(\{y\}) = b] \cap R(w))$ is defined, by regularity, and by [RP] it
 293 is equal to b . Hence $y \in R(w)$, and so wRy , as desired.

294 To show symmetry, assume xRy . Hence $y \in R(x)$, and so $P_x(\{y\}) = a > 0$.
 295 Thus $x \in [P(\{y\}) = a]$. Moreover, by reflexivity we have $x \in R(x)$. Hence
 296 $P_{prior}(\{y\} \cap [P(\{y\}) = a] \cap R(x)) / P_{prior}([P(\{y\}) = a] \cap R(x)) = P_x(\{y\} \mid$
 297 $[P(\{y\}) = a])$ is defined, by regularity, and by [RP] it is equal to a . Now suppose,
 298 for reductio, that $a < 1$. For transitivity, for every z , if $z \in R(y)$, then $R(z) \subseteq R(y)$.
 299 So $P_z(R(y)) = 1$. Then $[P(R(y)) = a] \cap R(y) = \emptyset$. So $P_{prior}(R(y) \cap [P(R(y)) =$
 300 $a] \cap R(x)) / P_{prior}([P(R(y)) = a] \cap R(x)) = 0$. Then, by [RP], $a = 0$. Contradiction.
 301 Hence $a = 1$. Thus $P_x(R(y)) = 1$. Thus we obtain that $x \in R(x) \subseteq R(y)$, so
 302 $x \in R(y)$, and hence yRx , as desired. \square

303 The upshot is that, given that we have not demanded that R be an equivalence
 304 relation, violations of [RP] can be true in some structure S . Moorean statements can
 305 be true as well, of course. However, doxastic Moorean statements cannot be themselves
 306 believed, while epistemic Moorean statements cannot themselves be known to be true
 307 (even if they are, as a matter of fact, true). By contrast, it is not difficult to see that if R
 308 is not an equivalence relation, violations of [RP] can be given maximum probability
 309 in S . So [RP Failure] is not the probabilistic equivalent of a Moorean statement.

310 Someone could object at this point that all we have shown is that R should be an
 311 equivalence relation, out of rationality considerations. I will address this objection
 312 with some detail in Sect. 5. But first, let us make sure that [RP Failure] can indeed
 313 have maximum evidential probability in S .

314 4 Longing for Fitch's paradox: paradox unduly lost

315 Evidently, if [RP Failure] is a genuine quasi-Moorean statement, and [RP Failure]
 316 is sometimes justified, then quasi-Moorean statements are sometimes justified, which
 317 presumably means that quasi-Moorean statements, unlike their non-probabilistic coun-
 318 terparts, can sometimes be legitimately asserted. Indeed, we can prove that, within the
 319 present setting, violations of particular instances of [RP] can in fact receive maximum
 320 probability (actually, they are knowable) and hence that quasi-Moorean statements
 321 so understood do not trigger a probabilistic version of Fitch's paradox: the modal
 322 collapse¹⁸ is suitably blocked. Therefore, either there is no probabilistic analogue of
 323 Fitch's paradox—at least not in a Kripke setting like the one presented here—or [RP
 324 Failure] is not a genuine quasi-Moorean statement.

325 Let $S^* = \langle W, R, R^*, P_{prior}, v \rangle$ be a structure where W, R, P_{prior} and v are as
 326 before, and where R^* is an alethic accessibility relation. Consider the following version
 327 of a Probabilistic Knowability Principle, formulated in the metalanguage:

328 $\forall \varphi \forall w$: if $\models_w \varphi$, then $\models_w \diamond [P([\varphi]) = 1]$ [P Knowability]

¹⁸ Or whatever it is that we think Fitch's paradox shows. Cf. the last paragraph of Sect. 1.

329 Or, equivalently:

330 $\forall\varphi\forall w$: if $\models_w \varphi$, then $\exists x wR^*x : P_x([\varphi]) = 1$ [P Knowability']

331 Consider also the following instance of [P Knowability], for $r \neq s$:

332 If $\models_x [P([\varphi] \mid [P([\varphi]) = r]) = s]$, then
 333 $\models_x \diamond[P([P([\varphi] \mid [P([\varphi]) = r]) = s]) = 1]$

334 Essentially, it says that, if a particular violation of [RP] is true in a given world,
 335 then it is possible for an agent to assign probability 1 to such a violation. It is easy
 336 to check that there are models in which both the antecedent and the consequent of
 337 this conditional come out true. All we need to do is find some proposition A such
 338 that $[P(A \mid [P(A) = r]) = s] \subseteq R(x)$, for $r \neq s$, for some world x that can reach to
 339 itself through R^* —which is a fairly reasonable and modest requirement to impose on
 340 alethic modalities. To illustrate this situation, let us just enrich our prior toy example
 341 from Fig. 2 with the assumption that wR^*w . As the intersection between A and
 342 $[P(A) = 1/2]$ is empty, we obtain:

343 $P_w(A \mid [P(A) = 1/2]) = 0$
 344 $P_y(A \mid [P(A) = 1/2]) = 0$
 345 $P_x(A \mid [P(A) = 1/2]) = 0$

346 Indeed, the set of worlds in which ' $P_w(A \mid [P(A) = 1/2])$ ' is 0 coin-
 347 cides with $R(w)$; hence $P_w([P(A \mid [P(A) = 1/2]) = 0]) = 1$, and,
 348 consequently, $\models_w [P([P(A \mid [P(A) = 1/2]) = 0]) = 1]$. Insofar as wR^*w ,
 349 ' $\diamond[P([P(A \mid [P(A) = 1/2]) = 0]) = 1]$ ' is true in w .¹⁹

350 As we can see, we obtained a true instance of [P Knowability] in which we assigned
 351 maximum probability to a true (potential) quasi-Moorean statement, so the probabilis-
 352 tic analogue to Fitch's paradox could not get off the ground. This should be disturbing.
 353 Paradoxes do not just dissolve in the air; we cannot get rid of a paradox without at
 354 least explaining what went wrong the first time we presented it. Regardless of what
 355 our favorite diagnosis concerning Fitch's original argument is, we should expect some
 356 basic structural features of knowledge attribution to be preserved in a more general set-
 357 ting. Notice that maximal evidential probability behaves as a knowledge-like concept,
 358 in the sense that, for a finite W , $P_w([\varphi]) = 1$ iff $\models_w K\varphi$, for any φ and w ²⁰ (if W is
 359 infinite, then ' $P_{prior}([\varphi] \cap R(w)) / P_{prior}(R(w))$ ' could be 1 even when $R(w) \not\subseteq [\varphi]$,
 360 so while the conditional 'if $\models_w K\varphi$, then $P_w([\varphi]) = 1$ ' is bound to be true, the
 361 converse need not hold).

362 It might be that we cannot do better within the present framework. Or, more likely,
 363 we can suspect that our candidate for a quasi-Moorean statement was not good enough.
 364 Let me put it as a dilemma: Either there is a serious, insurmountable asymmetry

¹⁹ As a matter of fact we have obtained something stronger, to wit, we have obtained that the sentence stating that [RP Failure] has probability 1 is actually true in w .

²⁰ Recall that $P_w([\varphi]) = P_{prior}([\varphi] \mid R(w))$ is always well defined, due to the regularity of P_{prior} .

365 between the attribution of knowledge and that of maximal evidential probability, or
 366 something has been lost in translation, so to speak.

367 5 Discussion: reflection, rationality and vagueness

368 The dilemma from the previous section could be rejected if we restrict acceptable
 369 models to those which adopt an equivalence accessibility relation, *i.e.*, to S5. I do not
 370 think we should demand such a restriction. Let me elaborate on this point.

371 Consider the following objection to the idea that violations of [RP] can be suc-
 372 cessfully known. Some authors have contended that doxastic versions of Moorean
 373 statements are not *logically* problematic, in the sense that a (deeply troubled) agent
 374 could well believe, *as a matter of fact*, that p and that he doesn't believe that p .²¹
 375 To the extent that this is the case, there might be room to develop models in which
 376 believing such abnormal statements is possible; we may well rely on a paraconsistent
 377 framework, or impose restrictions on the B operator, among other options. Analo-
 378 gously (so the objection goes) we may well find models—such as those presented in
 379 the previous section—in which violations of [RP] not only come out true, but where
 380 agents may also know that this is the case. However, this does not make such viola-
 381 tions any less troubling, in the same sense that building a model that makes it possible
 382 for agents to believe Moorean assertions does not dispel the air of paradox we feel
 383 in the original Moorean case. Rather, what we should say is that such models do not
 384 capture the behavior of *rational* attitudes. An ideally rational agent is one whose R
 385 is an equivalence relation, just as a rational agent is one who does not believe con-
 386 ceptual impossibilities. Under this perspective, a rational agent just *cannot* violate the
 387 Reflection Principle, let alone *know* that he has violated it. To put it differently, our
 388 fan of [RP] can object along these lines: trying to convince a supporter of [RP] that
 389 one can be rational and still know [RP Failure] is akin to trying to convince a classical
 390 logician that believing a doxastic Moorean statement is also not irrational, because we
 391 can always find consolation in paraconsistent logic.²²

392 To address this objection let me re-assess whether an equivalence accessibility
 393 relation is indeed mandatory for rational agents. As I can see it, we have good reasons
 394 to resist this claim. The reasons I have in mind differ from the usual complaint according
 395 to which an S5 model, though simpler and more tractable, depicts an unacceptably
 396 strong version of ideal rationality. In general, I am not a big fan of attempts to debunk
 397 a particular formal account on the grounds that it is too much idealized—idealizations
 398 can play an important role at the time of clarifying myriads of notions. Rather, the
 399 problem in this case is that S5 models give us *the wrong kind of idealization*. The
 400 problem is not that real agents typically do not verify, say, transparency claims (such
 401 as those embodied in the *KK* Principle) but, rather, that *ideal* agents should be sensitive
 402 to vagueness considerations in a way that S5 agents cannot be.

²¹ This is one of the reasons why Kvanvig (2006) claims that there is no interesting analogy between Fitch's and Moore's paradox.

²² Paraconsistent analyses of Fitch's paradox (such as Beall (2009)) argue that agents can indeed know epistemic Moorean statements; on similar grounds, a paraconsistent logician may well disagree with the claim that an agent just cannot, as a matter of logic, rationally believe a doxastic Moorean statement.

403 Consider worlds w_i , for $i = 1 \dots n$, such that the length of a particular table is 50
 404 + i cm in each world w_i . Let ‘ φ ’ be ‘the table is less than 100 cm long’. ‘ φ ’ is true in
 405 worlds w_1 to w_{49} . It can be intuitively appealing to assume that we can determine the
 406 length of the table with an error of ± 3 cm (of course the example can be modified to
 407 make it as realistic as we want, but for our current purposes this will suffice). In this
 408 scenario, transitivity is violated: in world w_{46} we know that φ , but we do not know
 409 that we know this, since in world w_{47} , which is epistemically accessible to w_{46} , we
 410 no longer know that φ . In general, demanding transitivity would make it impossible
 411 to account for cases in which we would be inclined to attribute knowledge ‘locally’,
 412 so to speak; this is one of Williamson’s points in Williamson (2000), ch.7. As is well
 413 known, Williamson also uses examples of this sort to explain the erosion that occurs
 414 when we go up to higher-order levels of knowledge. Our knowledge of ‘ φ ’ becomes
 415 more uncertain the closer we get to w_{50} , so at some point it becomes natural to say
 416 that we do not know that we know that φ , even if we in fact do know that φ : a naïve
 417 sense of epistemic transparency (related to the truth of the KK Principle) is bound
 418 to fail. Now, nothing in my example so far indicates the presence of vague terms at
 419 work;²³ however, we can change the story slightly so that vagueness becomes the main
 420 issue. Just take ‘ φ ’ to be ‘the table is large’, or change the scenario to fit your favorite
 421 example of vagueness. Let me bracket here the problem of how to assign truth values
 422 to such a sentence in intermediate worlds (as this would go beyond the purposes and
 423 goals of this paper). Still, it is clear that ideal agents should be discriminative enough
 424 so as to know what to say in extreme cases of application of the vague term.

425 One philosopher’s *modus ponens*, however, is another one’s *modus tollens*: A fan
 426 of S5 can object here that all this argument shows is that, in this and similar scenarios,
 427 ideal agents should discriminate better—so it should not be true that $(w_i, w_{i+1}) \in R$
 428 in the first place. But this will often lead us to say that R should be the identity relation,
 429 out of rationality considerations. However, it is not clear why we should assume that
 430 ideally rational agents are *empirically* omniscient beings, so demanding that R be the
 431 identity relation will be just inadequate for most cases.

432 There is a slightly different argument we can give here to reinforce this conclusion.
 433 It may well be that a possible explanation for the existence of vague terms in the
 434 language is that we often do not need more precise devices, given, among other things,
 435 that our discrimination powers are not perfect. This is related to the failure of empirical
 436 omniscience, rather than to a failure of reasoning capabilities. Insofar as we are trying
 437 to model perfect reasoners endowed with languages designed to encode imperfect
 438 discrimination capabilities, an identity accessibility relation is not the most appropriate
 439 tool, as it makes it impossible to use the resources of the language to their full potential.
 440 In short, a case can be made for the claim that it is the very existence of vague terms
 441 in the language which gives support to the convenience of treating ideal reasoners as
 442 exhibiting non-transitive accessibility relations.

443 An analogous phenomenon takes place within higher-order probabilities: here, too,
 444 agents may become increasingly more uncertain. Formally speaking, this is once again

²³ Actually, Williamson’s own diagnosis is that this particular phenomenon is not due to any putative vagueness related to the concept of knowledge. I agree; just to be clear, although I do think there is an interesting connection with vagueness here, it is not due to the vagueness of *knowledge*.

445 a straightforward consequence of allowing for non-transitive R s. Thus, in the light
 446 of the previous paragraph, there are no grounds for demanding that second order
 447 probabilities be always 0 or 1 *on pain of irrationality*. It is, again, a type of failure of
 448 self-knowledge – motivated, among other things, by sensitivity to vagueness-related
 449 considerations.

450 Here is another way of making the same point. There are two senses of reflection
 451 at stake: (i) the first one relates to self-knowledge, while (ii) the second one refers to
 452 a particular sense of probabilistic coherence, which is at the same time a potentially
 453 useful device to bridge the gap between prior and posterior probabilities. As it turns out,
 454 once sense (i) is elaborated in a way that makes room for the fact that ideal reasoners
 455 can be creatures endowed with vague languages (and less-than-perfect discrimination
 456 powers), it exerts immediate constraints on sense (ii). I will come back to these two
 457 senses of reflection in Sect. 8.

458 Incidentally, let me point out that these observations are compatible with the defense
 459 of a moderate version of epistemic transparency; to do so we may resort to models
 460 in which the behavior of higher order knowledge operators is governed by different
 461 accessibility relations R^i , for $i = 1, 2, \dots$. Then the lack of transitivity in our epistemic
 462 accessibility relations can still lead to (weaker) versions of the *KK* Principle (such as
 463 ‘ $K^i\varphi \rightarrow K^{i+1}K^i\varphi$ ’), provided R^i and R^{i+1} are related in an appropriate way.²⁴ Such
 464 models can also verify moderate versions of [RP]. I am not going to dig any deeper
 465 on this point here, though the topic will come up again in later sections.

466 Further considerations might speak in favor of the convenience of abandoning
 467 symmetry as well, at least for some scenarios. Hence, once again, it is not clear it is
 468 a rationality requirement. Suppose in w_1 Sasha is not feeling quite well, even though
 469 she is not running a fever; by contrast, she does have a fever in w_2 . In w_1 Sasha is in
 470 doubt as to whether the real world is w_1 or w_2 ; if she were in w_2 , by contrast, she will
 471 be certain of her having a fever (though in w_1 she does not know that she would be so
 472 certain, say, because she is not aware of the accessibility structure of this framework).

473 In short, there are good reasons not to demand that R be always an equivalence
 474 relation, out of rationality. In any case, for our present purposes a milder claim will
 475 suffice: we just need to agree that failing to demand such an R (and hence failing to
 476 demand that [RP] be satisfied) is not as hard to digest as acknowledging the possibility
 477 of rationally believing a (doxastic) Moorean statement. In other words, it does not
 478 amount to the same radical departure from usual pre-theoretical notions of what we
 479 should expect from ideal reasoners. I trust we can secure at least this basic agreement.

480 6 In search of a probabilistic version of the knowability paradox

481 As it happens, there *are* alternative formulations for quasi-Moorean statements within
 482 setting \mathcal{S} , which yield the probabilistic equivalent to structural unknowability and give
 483 rise to probabilistic versions of Fitch’s paradox—so the symmetry can be restored.
 484 Consider, as we did in Sect. 3, a statement such as ‘ $\varphi \wedge \underline{P([\varphi]) < 1}$ ’. It is easy to see

²⁴ See Sect. 9.

485 that, although this conjunction can be true in some worlds, its evidential probability
 486 can never be 1, in any world. In other words, we can prove that:

487 **Proposition 2** *Let S be as before. Then, $\forall w : P_w([\varphi] \cap [P([\varphi]) < 1]) < 1$.*

488 *Proof* Take any world $w \in W$. By definition, $P_w([\varphi] \cap [P([\varphi]) < 1]) = P_{prior}([\varphi] \cap [P([\varphi]) < 1] \cap R(w)) / P_{prior}(R(w))$. For this
 489 probability to be 1, proposition ‘ $R(w)$ ’ needs to be included both in ‘ $[\varphi]$ ’ and in
 490 ‘ $[P([\varphi]) < 1]$ ’. But, although the intersection between ‘ $[\varphi]$ ’ and ‘ $[P([\varphi]) < 1]$ ’ need
 491 not be empty (hence the conjunction of the relevant sentences can be true in w), and
 492 even though ‘ $R(w)$ ’ could in principle be included in any of the two, it cannot be
 493 included in both. If ‘ $R(w)$ ’ is included in ‘ $[\varphi]$ ’, then $P_w([\varphi]) = 1$, and hence, insofar
 494 as R is reflexive, $w \in [\varphi]$ (so $w \notin [P([\varphi]) < 1]$). On the other hand, if $R(w)$ is
 495 included in $[P([\varphi]) < 1]$, this means that all worlds in $R(w)$ can reach at least one
 496 not- φ world. So ‘ $\varphi \wedge [P([\varphi]) < 1]$ ’ cannot have maximal evidential probability: it is
 497 probabilistically unknowable. \square

499 Incidentally, notice that the truth of Proposition 2 is independent of how we choose
 500 the alethic relation R^* . Notice also that this inequality already incorporates the possi-
 501 bility to account for vague probabilities, to the extent that they can be cashed out by
 502 intervals $[r, s] \subseteq [0, 1]$.

503 Let us see now how a Fitch-like argument could go, using the quasi-Moorean
 504 statement just suggested, and a principle such as:

505
$$\forall \varphi \forall w : \text{if } \models_w \varphi, \text{ then } \exists x : P_x([\varphi]) = 1 \quad \text{[P Knowability 2]}$$

506 The proof will proceed in the metalanguage, insofar as expressions such as ‘ $P_w([\varphi]) =$
 507 r ’ are formulated by the theoretician. For any structure S :

- 508 1. $\exists w : P_w([\varphi] \cap [P([\varphi]) < 1]) = 1$ [assumption]²⁵
- 509 2. $\neg \exists w : P_w([\varphi] \cap [P([\varphi]) < 1]) = 1$ [by Proposition 2]
- 510 3. $\forall \varphi \forall w : \text{if } \models_w \varphi, \text{ then } \exists x : P_x([\varphi]) = 1$ [P Knowability 2]
- 511 4. If $\models_w (\varphi \wedge [P([\varphi]) < 1])$, [instance of [P Knowability 2]]
 512 then $\exists x P_x([\varphi] \cap [P([\varphi]) < 1]) = 1$
- 513 5. If $\models_w \varphi$ and $P_w([\varphi]) < 1$, [from 4]
 514 then $\exists x P_x([\varphi] \cap [P([\varphi]) < 1]) = 1$
- 515 6. If $\models_w \varphi$, then $P_w([\varphi]) = 1$ [from 5, 2]
- 516 7. $\forall \varphi \forall w : \text{if } \models_w \varphi, \text{ then } P_w([\varphi]) = 1$ [generalization from 6, given
 517 that ‘ φ ’ was any proposition
 whatsoever, and ‘ w ’ was also
 any world whatsoever.]
- 518 8. If for all worlds and sentences, if $\models_w \varphi$, [3 – 7]
 519 then $\exists x P_x([\varphi]) = 1$, then for all worlds and
 sentences, if $\models_w \varphi$, then $P_w([\varphi]) = 1$

²⁵ Notice that here I am trying to mimic standard proofs of Fitch’s result. Such proofs typically start by assuming that a Moorean statement can be known. Likewise, here I start by assuming that a quasi-Moorean statement (i.e., ‘ $\models_w (\varphi \wedge [P([\varphi]) < 1])$ ’) can receive maximum evidential probability

520 In other words, if all propositions that are true in a particular world are given
 521 maximal evidential probability in *some* world, then all propositions that are true in a
 522 world are given maximal evidential probability in that very same world. So R is bound
 523 to be the identity relation.

524 Notice that we have not required Necessitation in step (2), even though it is easy to
 525 show that (2) is equivalent to a statement with a modal operator:

$$526 \quad 2'. \forall x \neg \exists w x R^* w : \models_w P([\varphi] \cap [P([\varphi]) < 1]) = 1, \text{ i.e.:}$$

$$527 \quad 2''. \forall x \models_x \square P([\varphi] \cap [P([\varphi]) < 1]) < 1$$

528 Notice, moreover, that [P Knowability 2] is weaker than [P Knowability] as I used
 529 it in Sect. 4; there are no diamonds in step (3) of the proof as I have just presented it. [P
 530 Knowability 2] demands that every truth have maximal evidential probability in some
 531 world (rather than: in some of the worlds that relate alethically to a given world). So
 532 we need not consider an alethic R^* at all (or, equivalently, we can say that R^* is the
 533 universal relation). This simplifies the proof a bit, without loss of generality. Actually,
 534 as we relied on a weaker Probabilistic Knowability Principle, the result we obtained
 535 is stronger than the one we would have attained with the aid of [P Knowability]. In
 536 any case, notice that a similar simplification could have been applied to the original,
 537 non-probabilistic Fitch's paradox.

538 7 Generalizing the knowability principle

539 We can generalize what we have presented so far and consider an even weaker formu-
 540 lation for the Probabilistic Knowability Principle, along the following lines:

$$541 \quad \text{If } \models_w \varphi, \text{ then } \exists x : P_x([\varphi]) \geq r; \quad \text{[P Knowability 2']}$$

542 where r is a threshold for, say, 'the agent is confident enough' (given what she
 543 knows). Once again, we will obtain a Fitch-like paradox when ' φ ' is replaced by
 544 ' $\psi \wedge [P([\psi]) < r]$ '. To see this, notice that we can prove the following:

545 **Proposition 3** *Let S be as before. $\forall w$: If $w \in [P([\varphi]) < r]$, then $P_w([\varphi] \cap [P([\varphi]) < r]) < r$*

547 *Proof* Assume $w \in [P([\varphi]) < r]$. Then by definition $P_{prior}([\varphi] \cap R(w))/P_{prior}$
 548 $(R(w)) < r$. Hence, we also have $P_{prior}([\varphi] \cap [P([\varphi]) < r] \cap R(w))/P_{prior}(R(w)) =$
 549 $P_w([\varphi] \cap [P([\varphi]) < r]) < r$. \square

550 Now we can again obtain a Fitch-like claim to the effect that, if all truths are such
 551 that we could become confident of them, then all truths are such that we are currently
 552 confident that they are indeed true. As before, the proof proceeds in the metalanguage.
 553 Then, for any structure S :

- 554 1. $\forall \varphi \forall w$: if $\models_w \varphi$, then $\exists x : P_x([\varphi]) \geq r$ [[P Knowability 2'], for some
 555 threshold $r \neq 0$]
- 556 2. If $\models_w (\varphi \wedge [P([\varphi]) < r])$, [instance of [P Knowability 2']]
 557 then $\exists x P_x([\varphi] \cap [P([\varphi]) < r]) \geq r$

- 558 3. If $\models_w \varphi$ and $P_w([\varphi]) < r$, [from 2]
 559 then $\exists x P_x([\varphi] \cap [P([\varphi]) < r]) \geq r$
- 560 4. If $\models_w \varphi$ and $P_w([\varphi]) < r$, [by Proposition 3]
 561 then $\neg \exists x P_x([\varphi] \cap [P([\varphi]) < r]) \geq r$
- 562 5. If $\models_w \varphi$ and $P_w([\varphi]) < r$, then \perp [from 3 and 4]
- 563 6. If $\models_w \varphi$ then $P_w([\varphi]) \geq r$, [from 5]
- 564 7. If for all worlds and sentences, if $\models_w \varphi$, [1–6]
 565 then $\exists x P_x([\varphi]) \geq r$, then for all worlds and
 sentences, if $\models_w \varphi$, then $P_w([\varphi]) \geq r$

566 In short, potential confidence entails *actual* confidence. To put it in a somewhat
 567 different terminology, the ability to acquire confirmation collapses into actual con-
 568 firmation. Notice that here we are no longer dealing with probability 1—which is,
 569 arguably, a knowledge-like notion. So we are no longer dealing with a suitable transla-
 570 tion of Fitch’s result to a probabilistic realm, but with a genuine probabilistic perplexity,
 571 in Fitch’s spirit.

572 8 The reflection principle again: Van Fraassen’s integrity defense

573 In Sects. 3, 4 and 5 I argued that genuine quasi-Moorean statements cannot be identified
 574 with [RP Failure]. It might be contended, however, that the uneasiness we experience
 575 towards the violation of [RP] is part of the *explanation* of why certain *other* statements,
 576 such as those considered in Sects. 6 and 7, are genuine quasi-Moorean sentences.
 577 According to van Fraassen:

578 [L]et us note the formal connection, at least, between Moore’s paradox and
 579 the Reflection Principle. If we try to generalize Moore’s sentence schema to a
 580 probabilistic form, we arrive at:

581 It seems certain [likely, very likely] to me that: A and it seems unlikely to
 582 me that A ;

583 or, less qualitatively:

584 It seems likely to me to degree y that (A and it seems likely to me to degree
 585 x that A): $P(A \wedge p(A) = x) = y$

586 where the number x is lower than the number y and ‘ p ’ describes present (current)
 587 opinion. The synchronic form of the Special Reflection Principle ...is violated
 588 unless y is less than or equal to x . (van Fraassen (1995), p. 19)

589 Clearly, ‘ $P(A \mid P(A) = r) = r$ ’ entails that $P(A \cap p(A) = r) \leq r$. To adjust
 590 the terminology to our current framework, if r is less than 1, attributing maximum
 591 probability to ‘ $\varphi \wedge [P([\varphi]) = r]$ ’ entails that [RP] has been violated. Thus we might
 592 suggest that the paradoxicality of genuine quasi-Moorean statements lies in the fact
 593 that they yield violations of [RP]. However, *pace* van Fraassen, *we do not need [RP] to*
 594 *account for the Moorean blindspot*. As we have seen, ‘ $P_w([\varphi] \cap [P([\varphi]) < 1]) = 1$ ’
 595 is an impossible claim in S , for any w and φ (Proposition 2); [RP] does not play any
 596 role at the time of determining this impossibility. This does not mean to say that there

597 cannot be other reasons to defend [RP], of course—but a putative connection with
598 Moore's paradox does not seem to be one of them.

599 The last point can be contested on the following grounds. Ultimately, van Fraassen's
600 contention is that satisfying [RP] is crucial to our *integrity*. Moore's paradox and [RP
601 Failure] share in this respect the same type of 'inconsistency in a broad sense' which
602 is pragmatic rather than semantic:

603 It seems to me therefore that the correct notion of probabilistic incoherence must
604 take its inspiration from the notion of inconsistency made manifest by Moore's
605 paradox.... It is not inconsistent in the sense of 'unsatisfiable', 'incapable of
606 being true', which is the semantic notion of inconsistency. But I cannot have
607 a coherent state of opinion which I could express by a statement of the form
608 $[\varphi \wedge \neg B\varphi]$ (van Fraassen (1995), p. 27).

609 This is indeed an appealing defense, but the problem is that integrity can be inter-
610 preted in many ways. There is a different sense of integrity we might also feel pressed
611 to honor—and there is a potential conflict between the two. The tension appears more
612 forcefully once we take epistemic changes into account, as in diachronic versions of
613 the Reflection Principle.

614 Let us consider once again the distinction drawn in Sect. 5 between two senses of
615 reflection: Reflection as transparency (related to self knowledge), and reflection as a
616 type of (probabilistic) coherence. These two senses of reflection relate in turn to two
617 senses of integrity. According to the first one, integrity requires that we take pride in
618 feeling accountable for our present and future actions; they are all (robustly) ours. It
619 is precisely because we feel we ought to be so accountable that we better know who
620 we are—where who we might become is also part of who we are.²⁶ It might happen
621 that we do not approve of the way in which we foresee we will change our minds,
622 in which case we might be able to take measures to alleviate the mistake (from our
623 current point of view) as Ulysses did when he imagined himself facing the sirens.
624 Identity, and responsibility within identity, overrides coherence over time. Let me call
625 it 'Ulysses integrity'. Ulysses-type integrity can account for a temperate version of the
626 *KK* Principle, as well as for temperate versions of [RP],²⁷ including also diachronic
627 counterparts of such temperate versions, but not for the full-blown [RP]—not even for
628 the synchronic case, as we have seen.

629 A second sense of integrity, by contrast, requires that we take pride in coherence. If
630 our future self, as we foresee it, is not quite the person we believe we should become,
631 we give up on such a future person. Diachronic coherence overrides personal identity
632 over time (and notice that the thesis works both ways: the present self could give up
633 on his or her past self). It is the sense of integrity that leads the idealist young man to
634 say to his lover: "if in the future I abandon my ideals, I beg you to think that the person
635 you now know and love does no longer exist; in such event *I* will be dead". Let me

²⁶ Of course, we cannot demand knowledge of the future on rationality grounds. What can be demanded, however, is that we take active steps to be able to make accurate predictions about our future temporal slices.

²⁷ Cf. Cresto (2012).

636 call it ‘Parfit integrity’.²⁸ A consequence of Parfit integrity is that I consider a future
 637 person to be identical with me only to the extent that I can have trust in that future
 638 person’s probability assignment. Parfit-type integrity can then account for diachronic
 639 versions of [RP], and, by extension, for their synchronic, more restricted versions.

640 In short, according to van Fraassen Moore’s paradox exhibits a type of ‘inconsistency
 641 in a broad sense’ which is also the type of inconsistency we can identify in some
 642 cases of integrity failure—what I have called ‘Parfit integrity’. However, there are
 643 other types of integrity failure we might worry about. Since Parfit integrity is not all
 644 we demand from agents, it is not clear whether [RP] is a principle we should always
 645 enforce; at any rate, not out of Moorean-type considerations.

646 9 Beyond epistemic paradoxes. Motivations for hierarchic languages

647 One of the main goals of this paper was to explore the structure of successful
 648 quasi-Moorean statements, and, concomitantly, to build a probabilistic version of the
 649 Knowability Paradox. The results obtained helped us draw a number of morals that go
 650 well beyond Fitch’s argument. We have shown that violations of [RP] can be attributed
 651 maximal evidential probability, so [RP Failure] is not itself a quasi-Moorean sentence:
 652 the Moorean spirit has been ‘lost in translation’. More generally, the ‘integrity defense’,
 653 which seeks to compare violations of [RP] with Moorean-type irrationality, should be
 654 taken *cum grano salis*, since [RP] may actually conflict with other senses of integrity.
 655 Of course, there might well be alternative strategies to show the putative illegitimacy
 656 of [RP Failure]—Dutch Books, calibration arguments, etc. But it should be clear by
 657 now that at least some lines of argument will not do.

658 In addition, we have confirmed that the chosen formal framework maintains a
 659 healthy symmetry between knowledge attribution and the attribution of maximal evi-
 660 dential probability. Of course, symmetry as I understand it here is just a *necessary*
 661 condition for a satisfactory framework, but it is not sufficient. Moreover, if we take the
 662 symmetry criterion seriously, any adequate *solution* to the paradox should be unified as
 663 well. In order to do so, a more sophisticated setting might be desirable. In what follows
 664 I will outline such a setting, without entering into the details. As we will see, it leads
 665 naturally to a unified solution to both strands of Fitch’s paradox. Given the symmetry
 666 criterion, this fact can be turned into indirect evidence in favor of the formalism.

667 The system I favor is largely based upon the one we have been working with so far,
 668 but assumes a hierarchy of *K* and *P* operators, together with certain restrictive rules
 669 for building well-formed formulas. The motivation for this proposal goes as follows.

670 To begin with, a case can be made for the claim that second order probabilities
 671 demand that we conditionalize on second order evidence. Suppose we have information
 672 about the state of the weather tomorrow. We have read the forecast in the newspaper,
 673 watched the weather channel, etc. On the basis of all this information, we conclude
 674 that the probability of rain tomorrow in our city is 0.3. Now suppose a friend asks us

²⁸ See Parfit (1973), pp. 145–6. I do not mean to say that Parfit himself supports what I dubbed ‘Parfit Integrity’, but just that his description of what he calls ‘The Complex View’ comes close to capturing what I have in mind. Parfit’s example is re-elaborated in Elster (1984), Part II.

675 how probable it is that our rational degree of belief that there is rain tomorrow is in
 676 fact 0.3. As I see it, in this case our friend is no longer interested in the probability
 677 of a proposition about meteorology, but in the probability of a proposition *about the*
 678 *degree of confirmation* possessed by our original meteorological statement. Which is
 679 the relevant evidence to answer this question, then? Intuitively, what we have to assess
 680 is how good we are at the time of engaging in confirmation theory. Thus the relevant
 681 total evidence is no longer $R(w)$: the evidence for our second-order probability should
 682 consist in what we know about our capabilities to adequately confirm propositions;
 683 the strongest proposition that expresses this idea is in fact $KR(w)$.²⁹ Hence when we
 684 calculate a second order evidential probability we should conditionalize on $KR(w)$.
 685 The proposal then generalizes to increasingly higher levels.³⁰

686 Consider now a metalinguistic statement such as ' $P_w([P([\varphi]) = r])$ '. As we have
 687 seen, this statement is meant to capture the intuition that we are calculating a *second*
 688 order probability. However, the argument of the probability function is just a set of
 689 worlds; we could have well referred to it by other means—for example, by means of
 690 a suitable sentence of L without epistemic operators (say, ' $[\psi]$ '), as in a regular first
 691 order probability. Therefore, propositions understood as sets of worlds seem to be too
 692 coarse grained for what we want.

693 A possible suggestion at this point is to let the arguments of our probability functions
 694 be *sentences* of well-regimented languages; we can define a sequence of languages
 695 $L^0, L^1 \dots L^n \dots$, with probability operators $P^0, P^1 \dots P^n \dots$ that apply to sentences
 696 of lower lever languages. Thus, the probability of a set of worlds will depend crucially
 697 on the way we refer to it. In other words, at the time of calculating evidential
 698 probabilities, the 'mode of presentation' matters. For structural reasons of internal
 699 coherence, we should also demand a corresponding sequence of knowledge operators
 700 $K^0, K^1 \dots K^n \dots$, each with its own accessibility relation.³¹ It can be shown that, if
 701 we choose our accessibility relations carefully—so as to make sure that knowledge
 702 and probability attributions cohere with each other—the model can validate a moderate
 703 version of the *KK Principle*: to wit, if the agent has first order knowledge that
 704 φ , then she has *second order knowledge* that she has first order knowledge that φ
 705 (' $K^1\varphi \rightarrow K^2K^1\varphi$ '). Such a moderate version of *KK* can be satisfied without actually
 706 demanding that any of the accessibility relations be transitive. Hence the model can

²⁹ By definition, $KR(w) = \{x \in W : \text{if } xRy, \text{ then } y \in R(w), \text{ for all } y\}$. Hence $KR(w) \subseteq R(w)$.

³⁰ It might be contended that agents need not be aware of ' $KR(w)$ '—in which case they would not know which proposition they should conditionalize on (thanks to an anonymous referee for giving me the opportunity to clarify this point.). However, if there is a problem here, it is not exclusive of the enhanced framework, as similar considerations can be made for the standard setting; to wit, it might be contended that agents need not be aware of ' $R(w)$ ' either. There are at least two ways out, which relate to two very different interpretations of the formalism. On one hand, we can conceive of the framework as a tool for the theoretician (or the interpreter), who seeks to make knowledge and probability attributions from a third person point of view. She is the one who assesses, to the best of her knowledge, what the agent knows or ignores in each possible situation. On the other hand, we can think of the framework as 'viewed' from the inside, as it were, *i.e.*, as structured from the first person perspective. In this case we can take ' $R(w)$ ' to refer to the information the agent has consciously gathered. ' $KR(w)$ ' could then capture the subset of $R(w)$ which the agent takes to be the result of extremely reliable research methods, among other possibilities.

³¹ In a nutshell, by having the right sequence of knowledge operators we guarantee that statements with evidential probability 1 will be known by the agent.

707 still be sensitive to vagueness related considerations, as discussed in Sect. 5. All we
708 need is the weaker assumption that relation R^{i+1} composed with R^i is included in R^i .

709 Within this enriched setting we are no longer able to express Moorean or quasi-
710 Moorean statements in our sequence of languages.³² Expressions such as ‘ $(*)K^2(\varphi \wedge$
711 $\neg K^1\varphi)$ ’, as well as their probabilistic counterparts, are not well-formed formulas to
712 begin with. Therefore, the two versions of the Knowability Paradox dissolve. Unlike
713 other attempts to address Fitch’s paradox, in this case the syntactic restrictions on well-
714 formed formulas respond to principled reasons that are completely independent of the
715 discussion on Knowability. Once the restrictions are in place, however, we obtain a
716 unified answer to Fitch-like paradoxes as a nice side effect.

717 To sum up, the fact that we were able to develop a probabilistic Fitch-like argument
718 in the simpler setting tells us that it was a setting worth exploring; the fact that we could
719 not find straightforward solutions to the paradoxes within that framework suggests that
720 a more sophisticated structure might be welcome. Finally, the fact that we can obtain
721 a unified solution to the paradox within a richer framework, which has been originally
722 developed for independent reasons, reinforces the thought that we are on the right
723 track.

724 10 Concluding remarks

725 Along these pages I presented genuine examples of quasi-Moorean statements, and I
726 used them to build a probabilistic version of the Knowability Paradox. I also showed
727 that violations of [RP] do not share the relevant traits of Moorean-type irrationality:
728 Moorean and quasi-Moorean statements are true blind spots, whereas [RP Failure] is
729 not. Finally, I outlined a formal framework within which both strands of the paradox
730 received a uniform solution.

731 Let me address a few final concerns on the very structure and goals of this paper. As I
732 have already anticipated in the Introduction, someone might object that I have assumed
733 all along that the relevance of Fitch’s paradox lies in the fact that it reveals a worrisome
734 modal collapse between possible and actual knowledge,³³ and, analogously, between
735 possible and actual evidential probability. But it is far from clear whether this is the
736 right description of the problem.

737 The objection, however, would be misguided. All I have required is an endorse-
738 ment of what I dubbed ‘the symmetry criterion’ between knowledge and probability.
739 Whatever it is that we deem problematic with Fitch’s result, we should obtain a sim-
740 ilarly problematic result within a probabilistic framework. Hence, *if* Fitch’s original
741 argument involves a modal collapse of some sort, there should be an analogous modal
742 collapse within a probabilistic version of the argument, with the aid of a quasi-Moorean
743 statement. But this paper is neutral concerning the existence of such a modal collapse.

³² The proposed solution shares a family resemblance with other attempts to solve Fitch’s paradox with the aid of typed languages, such as Linsky (1986), Linsky (2009), or Paseau (2008). However, the sequence of languages that I have in mind is more restrictive than usual hierarchic proposals, in the sense that ‘ K^i ’ is only meant to apply to sentences of the form ‘ $K^{i-1}\varphi$ ’ or their negations; analogous restrictions apply to higher order probability statements. A rationale for this demand can be found in Cresto (2012).

³³ As suggested by Kvanvig (2006).

744 Actually, although my own diagnosis is that all versions of Moore's and Fitch's
 745 Paradox arise out of a confusion between levels of operators, the probabilistic version
 746 that I have offered here is compatible with many different interpretations of the source
 747 of the problem. Thus, for example, my probabilistic reconstruction of the Knowability
 748 argument is perfectly compatible with the claim that true verificationists should switch
 749 to intuitionistic logic in order to block the potential *reductio* of the Knowability Prin-
 750 ciple.³⁴ Alternatively, if we are convinced that there are good reasons to demand that
 751 the Knowability Principle be restricted to Cartesian Propositions,³⁵ we should equally
 752 demand that [P Knowability 2] be only instantiated by sentences that can be assigned
 753 probability 1 without contradiction.

754 On the other hand, it is not clear how a request to restrict the Knowability Principle
 755 to actual truths, as in Dorothy Edgington's proposal,³⁶ could work in the probabilistic
 756 scenario of Sects. 2, 3, 4, 5, 6, 7 and 8. Consider the following amendment to [P
 757 Knowability 2],³⁷ as a possible way of capturing the idea that only actual truths are
 758 knowable:

759 $\forall\varphi\forall w : \text{if } \models_w \varphi, \text{ then } \exists x : P_x(\{w\}) > 0, \text{ and } P_x([\varphi]) = 1$ [P Knowability A]

760 (or, equivalently, $\forall\varphi\forall w : \text{if } \models_w \varphi, \text{ then } \exists x : w \in R(x) \subseteq [\varphi]$). We can see
 761 that the paradox still runs:

762 4'. If $\models_w (\varphi \wedge [P([\varphi]) < 1])$, then [Instance of [P Knowability A]]
 763 $\exists x : P_x(\{w\}) > 0 \text{ and } P_x([\varphi] \cap [P([\varphi]) < 1]) = 1$

764 As ' $P_x([\varphi] \cap [P([\varphi]) < 1]) = 1$ ' is an impossible claim, we still obtain:

765 5'. If $\models_w \varphi$, then $P_w([\varphi]) = 1$

766 In other words, if there is no world x such that ' $[\varphi] \cap [P([\varphi]) < 1]$ ' can be included
 767 in $R(x)$, then the same is true for whichever world we can pick out in $R(x)$ as the
 768 actual world. So the rest of the argument follows without changes. We can take this
 769 result either as evidence that Edgington's amendment is not effective in this sort of
 770 probabilistic setting, or as a motivation to find a different translation for Edgington's
 771 intuition. I am not going to take a stance on this point here.

772 In any case, I submit that, once we adopt the symmetry criterion, finding a unified
 773 solution for both strands of Fitch's paradox is mandatory, if some solution is offered
 774 at all. In this sense, the symmetry criterion can be used not only to test the adequacy of
 775 a given probabilistic setting, but also to test the adequacy of different solutions to the
 776 knowledge version of the paradox. By itself, the criterion will not succeed in singling
 777 out a best answer, though some proposals can be ruled out as formally inadequate.
 778 This leaves us the interesting task of revising well known proposals in a systematic
 779 manner to see whether they pass the test. For the moment, we have learnt that there
 780 are formal refinements of Kripke settings that work just fine.

34 Williamson (1982), Dummett (2009).

35 Cf. Tennant (1997).

36 Edgington (1985), Edgington (2010). See also Rabinowicz and Segerberg (1994).

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